

A Glossary for IWGS (Auto-Generated)

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Preface

This document contains an English glossary for the course *Informatische Werkzeuge in den Geistes- und Sozialwissenschaften* at FAU Erlangen-Nürnberg (IWGS). It is automatically generated from the sources of the IWGS course notes and should be up-to-date with the course progress.

The glossary contains definitions for all technical terms used in the course, both the ones defined in the course, as well as the ones presupposed. The latter should be relatively few, since IWGS is intended as a beginner's course.

1 Glossary for IWGS

To make the role of arguments extremely clear, we write functions in **λ -notation**. For $f = \{(x, E) \mid x \in X\}$, where E is an expression, we write $\lambda x \in X.E$.

n -dim Cartesian space **n -dim Cartesian space:** $A^n := \{\langle a_1, \dots, a_n \rangle \mid 1 \leq i \leq n \Rightarrow a_i \in A\}$, call $\langle a_1, \dots, a_n \rangle$ a **vector**

n -dimensional Cartesian space An n -dimensional **Cartesian product** $A_1 \times \dots \times A_n$ is called a **n -dimensional Cartesian space** over A (and denoted A^n) iff $A_i = A$ for some set A for all i . We call $\langle a_1, \dots, a_n \rangle \in A^n$ a **vector**.

n -fold Cartesian product Let $A := \{A_i \mid 1 \leq i \leq n\}$ be a collection of sets, then the **n -fold Cartesian product** $A_1 \times \dots \times A_n$ is $\{\langle a_1, \dots, a_n \rangle \mid a_i \in A_i \text{ for all } 1 \leq i \leq n\}$, we call $\langle a_1, \dots, a_n \rangle \in A_1 \times \dots \times A_n$ an **n -tuple**. n is called the **dimension** of $A_1 \times \dots \times A_n$.

n -fold Cartesian product **n -fold Cartesian product:** $A_1 \times \dots \times A_n := \{\langle a_1, \dots, a_n \rangle \mid \forall i. 1 \leq i \leq n \Rightarrow a_i \in A_i\}$, call $\langle a_1, \dots, a_n \rangle$ an **n -tuple**

n -fold composition We write the **n -fold composition** of the relation R as R^n and define it by $R^1 := R$ and $R^{i+1} := \{S \circ R \mid S \in R^i\}$

n -tuple Defined along with **n -fold Cartesian product**

n -tuple Defined along with **n -fold Cartesian product**

p -closure Let p be a properties and $R \subseteq A \times B$ a **relation**, then we call the smallest (in terms of the \subseteq) relation $R' \supseteq R$ that has property p the **p -closure** of R .

p -closure Let p be one of the properties above and R be a relation, then we call the smallest relation $\supseteq R'R$ (in terms of the \subseteq) that has property p the **p -closure** of R .

ADT See **abstract data type**

AI See **Artificial Intelligence**

AI Defined along with **Artificial Intelligence**

AI See **Artificial Intelligence**

ALU Defined along with **central processing unit**

API See **application programming interface**

Information Interchange The **American Standard Code for Information Interchange** (ASCII) is a **character code** that assigns characters to numbers 0-127

Code	...0	...1	...2	...3	...4	...5	...6	...7	...8	...9	...A	...B	...C	...D	...E	...F
0...	NUL	SOH	STX	ETX	EOT	ENQ	ACK	BEL	BS	HT	LF	VT	FF	CR	SO	SI
1...	DLE	DC1	DC2	DC3	DC4	NAK	SYN	ETB	CAN	EM	SUB	ESC	FS	GS	RS	US
2...		!	"	#	\$	%	&	'	()	*	+	,	-	.	/
3...	0	1	2	3	4	5	6	7	8	9	:	;	<	=	>	?
4...	@	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O
5...	P	Q	R	S	T	U	V	W	X	Y	Z	[\]	^	_
6...	'	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o
7...	p	q	r	s	t	u	v	w	x	y	z	{		}	~	DEL

Artificial Intelligence **Artificial Intelligence (AI)** is intelligence exhibited by machines

Artificial Intelligence [Artificial Intelligence \(AI\)](#) is a sub-field of Computer Science that is concerned with the automation of intelligent behavior.

Artificial Intelligence [Artificial Intelligence \(AI\)](#) studies how we can make the computer do things that humans can still do better at the moment.

Boolean Defined along with [integer](#)

CLI See [command-line interface](#)

CPU See [central processing unit](#)

CS See [computer science](#)

Cartesian product The [Cartesian product](#) of an arbitrary (possibly infinite) indexed family of sets is defined as $\prod_{i \in I} X_i := \{f: I \rightarrow \bigcup_{i \in I} X_i \mid f(i) \in X_i\}$.

Cartesian product [Cartesian product](#): $A \times B := \{(a, b) \mid a \in A \wedge b \in B\}$, call (a, b) [pair](#).

Content MathML Defined along with [Mathematics Markup Language](#)

DAG We will sometimes use the abbreviation [DAG](#) for “directed [acyclic](#) graph”.

DAG Defined along with [cyclic](#)

DELETE Defined along with [method](#)

DOM See [document object model](#)

DTM Defined along with [nondeterministic Turing machine](#)

Extensible Markup Language See [XML](#)

FLOSS See [Free/Libre/Open-Source Software](#)

Open-Source Software [Free/Libre/Open-Source Software \(FLOSS\)](#) or just [open source](#) is software that is and licensed via [licenses](#) that ensure that its source code is available.

GET Defined along with [method](#)

Gregorian calendar The [Gregorian calendar](#) is a [solar calendar](#) that takes the birth of Christ is taken as the year 1 and has a complicated rule for leap years: Every year that is exactly divisible by four is a leap year, except for years that are exactly divisible by 100; the centurial years that are exactly divisible by 400 are still leap years.

HTML See [HyperText Markup Language](#)

HTML5 See [HyperText Markup Language](#)

HTTP See [Hypertext Transfer Protocol](#)

HTTP request Defined along with [user agent](#)

HyperText Markup Language The [HyperText Markup Language \(HTML\)](#), is a representation format for [web pages](#) [Hic+14].

HyperText Markup Language The [HyperText Markup Language \(HTML5\)](#), is believed to be the next generation of [HTML](#). It is defined by the W3C and the WhatWG.

Hypertext Transfer Protocol The [Hypertext Transfer Protocol \(HTTP\)](#) is an application layer protocol for distributed, collaborative, hypermedia information systems.

IDE See [integrated development environment](#)

ISO-Latin 16 Extensions of ASCII to 8-bit (256 characters) **ISO-Latin** 1 $\hat{=}$ “Western European”, ISO-Latin 6 $\hat{=}$ “Arabic”, ISO-Latin 7 $\hat{=}$ “Greek”...

International System of Units The **International System of Units** (SI) is the modern form of the metric system and is generally a system of [units of measurement](#) devised around seven **base units** and corresponding **dimensions**.

Internet The **Internet** is a worldwide computer network that connects hundreds of thousands of smaller networks. (The mother of all networks)

Kleene closure The operation of passing from an alphabet A to A^* is called **Kleene closure**, **Kleene operation**, or **Kleene star**. The operation \cdot^+ is called **Kleene plus**.

Kleene operation See [Kleene closure](#)

Kleene plus See [Kleene closure](#)

Kleene star See [Kleene closure](#)

Landau set The three **Landau sets** $\mathcal{O}(g), \Omega(g), \Theta(g)$ are defined as

- $\mathcal{O}(g) = \{f \mid \exists k > 0. f \leq_a k \cdot g\}$
- $\Omega(g) = \{f \mid \exists k > 0. f \geq_a k \cdot g\}$
- $\Theta(g) = \mathcal{O}(g) \cap \Omega(g)$

Landau set Let $g: \mathbb{N} \rightarrow \mathbb{N}$, then we define the three **Landau sets** $\mathcal{O}(g), \Omega(g), \Theta(g)$ as

- $\mathcal{O}(g) := \{g: \mathbb{N} \rightarrow \mathbb{N} \mid \exists k > 0. f \leq_a k \cdot g\}$
- $\Omega(g) := \{g: \mathbb{N} \rightarrow \mathbb{N} \mid \exists k > 0. f \geq_a k \cdot g\}$
- $\Theta(g) := \mathcal{O}(g) \cap \Omega(g)$

If $G \in \{\mathcal{O}(g), \Omega(g), \Theta(g)\}$, we often say that $f: \mathbb{N} \rightarrow \mathbb{N}$ is (of complexity) G , iff $f \in G$, accordingly we often write $f = G$ by a certain abuse of notation.

Given a particular function, e.g. $g: \mathbb{N} \rightarrow \mathbb{N}; n \mapsto n^3$, we often write $\mathcal{O}(g)$ as $\mathcal{O}(n^3)$, and analogously for $\Omega(g)$ and $\Theta(g)$.

MathTalk Abbreviations for Mathematical statements in **MathTalk**

- \wedge and \vee are common notations for *and* and *or*
- *not* is in mathematical statements often denoted with \neg
- $\forall x.P$ ($\forall x \in S.P$) stands for *condition P holds for all x (in S)*
- $\exists x.P$ ($\exists x \in S.P$) stands for *there exists an x (in S) such that proposition P holds*
- $\nexists x.P$ ($\nexists x \in S.P$) stands for *there exists no x (in S) such that proposition P holds*
- $\exists^1 x.P$ ($\exists^1 x \in S.P$) stands for *there exists one and only one x (in S) such that proposition P holds*
- **iff** as abbreviation for *if and only if*, symbolized by \Leftrightarrow
- the symbol \Rightarrow is used as a shortcut for *implies*

Mathematics Markup Language The **Mathematics Markup Language** (MathML) is an integrated framework for content and **presentation** in web-based mathematics.

Presentation MathML covers the basic font, box, grouping primitives for presenting the two-dimensional layout of mathematical formulae

Content MathML provides the an infrastructure for marking up the functional/logical structure of formulae as applications, variables, constants, and binding expressions.

NTM See [nondeterministic Turing machine](#)

ODF Defined along with [Office Open XML](#)

OOP See [object-oriented programming](#)

OOXML See [Office Open XML](#)

OS See [operating system](#)

Office Open XML Popular [word processors](#) include

- MS Word, an elaborated [word processor](#) for Windows, whose native format is **Office Open XML** (**OOXML**; file extension `.docx`).
- OpenOffice and LibreOffice are similar [word processors](#) using the **ODF** format (**Open Office Format**; file extension `.odf`) natively, but can also import other formats..
- Pages, a [word processors](#) for Mac OS X it uses a proprietary format.
- Office Online and GoogleDocs are browser-based real-time collaborative [word processors](#).

Open Office Format Defined along with [Office Open XML](#)

PDF See [Portable document format](#)

PGS work Defined along with [copyrightable work](#)

POST Defined along with [method](#)

PUT Defined along with [method](#)

Peano axioms The following set of axioms are called the **Peano axioms** ([Giuseppe Peano *1858, †1932](#))

Presentation MathML Defined along with [Mathematics Markup Language](#)

Portable document format **Portable document format** (**PDF**) is a [document format](#) that mixes text and graphics with a variety of content including logical structuring elements, interactive elements such as annotations and form-fields, layers, rich media (including video content), and three-dimensional objects using U3D or PRC, and various other data formats. The **PDF** specification also provides for encryption and digital signatures, file attachments, and metadata to enable workflows requiring these features.

REPL See [read-eval-print loop](#)

RTFM **RTFM** ($\hat{=}$ “read those [fine](#) manuals”)

RTFM **RTFM** ($\hat{=}$ “read the fine manuals”)

SI See [International System of Units](#)

SVG See [Scalable Vector Graphics](#)

Scalable Vector Graphics [Scalable Vector Graphics \(SVG\)](#) is an [XML](#)-based [markup](#) format for [vector graphics](#).

Turing complete An [information processing system](#) is said to be [Turing complete](#) or [computationally universal](#) if it can be used to simulate any [Turing machine](#).

Turing machine A [Turing machine](#) consists of

- An [tape](#) with infinitely many [cells](#), each of which contains a symbol from a finite alphabet \mathcal{A} with $\#(\mathcal{A}) \geq 2$.
- A [head](#) that can read/write symbols on the tape and move left/right.
- A [state register](#) that stores the state of the Turing machine. The set of states is finite, the [state register](#) is initialized with a special start state.
- An [action table](#) (or [program](#)) that – given the symbol it has just read from the tape and the state it is currently in – specifies the next [action](#), i.e. what symbol to write, how to move the head and the next state. If no entry applicable the machine [halts](#).

UCS See [universal character set](#)

UI See [user interface](#)

URI See [uniform resource identifier](#)

URL A [uniform resource locator](#) ([URL](#)) is a [URI](#) that that gives access to a [web resource](#), by specifying an access method or location. All other [URIs](#) are called [uniform resource names](#) ([URN](#)).

URL See [uniform resource locator](#)

URN Defined along with [uniform resource locator](#)

UTF-8 The [UTF-8](#) encoding encodes each character in one to four octets (8-bit bytes):

1. One byte is needed to encode the 128 US-ASCII characters (Unicode range U+0000 to U+007F).
2. Two bytes are needed for Latin letters with diacritics and for characters from Greek, Cyrillic, Armenian, Hebrew, Arabic, Syriac and Thaana alphabets (Unicode range U+0080 to U+07FF).
3. Three bytes are needed for the rest of the Basic Multilingual Plane (which contains virtually all characters in common use).
4. Four bytes are needed for characters in the other planes of Unicode, which are rarely used in practice.

WFH See [work made for hire](#)

WWW See [World Wide Web](#)

WWWWeb See [World Wide Web](#)

World Wide Web The [World Wide Web](#) ([WWW](#) or [WWWWeb](#)) is an open source information space where documents and other web resources are identified by [URLs](#), interlinked by hypertext links, and can be accessed via the [Internet](#).

XHTML [XHTML](#) is the [XML](#) version of [HTML](#). (just make it valid XML)

XML [XML](#) (short for [Extensible Markup Language](#)) is a framework for [markup formats](#) for [electronic documents](#) and structured [data](#).

XML document tree The **XML document tree** is made up of **element nodes**, **attribute nodes**, **text nodes** (and **namespace declarations**, **comments**,...))

XML literal We call any **string** that is well-formed **XML** an **XML literal**.

XML namespace An **XML namespace** is a string that identifies an **XML** vocabulary. Every element and attribute name in **XML** consists of a **local name** and a **namespace**.

XML path language The **XML path language (XPath)** is a language framework for specifying fragments of **XML** trees.

XPath See **XML path language**

absolute value The **absolute value** $|r|$ of an **integer** r is defined as $\begin{cases} r & \text{if } r \geq 0 \\ -(r) & \text{else} \end{cases}$.

absolute value The **absolute value** $|r|$ of an **rational number** $\frac{a}{b}$ is defined as $\frac{|a|}{|b|}$.

absolute value The **absolute value** $|r|$ of a number r is defined as $\begin{cases} r & \text{if } r \geq 0 \\ -(r) & \text{else} \end{cases}$.

absolute value Real **absolute value**, **addition**, **subtraction**, **multiplication**, **division**, and **exponentiation**, **square roots**, r -th **roots**, and the ordering relations are extended to the **real numbers**, so that they respect limits.

abstract data type An **abstract data type (ADT)** is a mathematical model for a **data types**, which specifies a **container** by its behavior from the point of view of a user of the data, specifically in terms of possible **value**, possible **operations** on **data** of this **type**, and the behavior of these **operations**.

accepted unit Some units that have important contemporary application worldwide, or are otherwise commonly encountered worldwide and can be expressed as scalar multiples of **derived units**. They have been given the status of **accepted units** in the **SI** system.

accepting state Defined along with **nondeterministic Turing machine**

action Defined along with **Turing machine**

action table Defined along with **Turing machine**

acyclic Defined along with **cyclic**

acyclic Defined along with **cyclic**

addition **Addition** is extended to the **integers** by defining the **sum** as

$$a + b := \begin{cases} |a| + |b| & \text{if } a, b \in \mathbb{N} \\ -(|a| + |b|) & \text{if } a, b \in \mathbb{Z}^- \\ |a| + (-(|b|)) & \text{if } (a \geq b) \\ |b| + (-(|a|)) & \text{if } (a < b) \end{cases}$$

addition We define **addition** on the **rational numbers**: the **sum** $\frac{a}{b} + \frac{c}{d}$ is $\frac{a \cdot d + b \cdot c}{b \cdot d}$.

addition **Addition** + computes the **sum** of $a + b$ of **natural numbers** a and b . It defined by the equations $x + 0 = x$ and $x + s(y) = s(x + y)$, where s is the successor function.

addition **Addition** + computes the **sum** of $a + b$ of **numbers** a and b .

addition Defined along with [absolute value](#)

addition operation We “define” the [addition operation](#) \oplus procedurally (by an algorithm)

- adding zero to a number does not change it.
written as an equation: $n \oplus 0 = n$
- adding m to the successor of n yields the successor of $m \oplus n$.
written as an equation: $m \oplus s(n) = s(m \oplus n)$

algorithm An [algorithm](#) is a formal or informal specification for solving a problem by executing a finite sequence of [instructions](#) (concrete or imaginary/abstract) [information processing systems](#).

alphabet Defined along with [formal language](#)

alphabet An [alphabet](#) A is a finite set; we call each element $a \in A$ a [character](#), and an n -tuple of $s \in A^n$ a [word](#) (or [string](#)) of over A . We will often write a [string](#) $\langle c_1, \dots, c_n \rangle$ as " $c_1 \dots c_n$ " or even as $c_1 \dots c_n$. We write the [empty word](#) ([empty string](#)) in A^0 with ϵ .

alphabet Defined along with [nondeterministic Turing machine](#)

amount Defined along with [length](#)

ancestor The [ancestor](#) and [descendant relations](#) are the [transitive closures](#) of the [parent](#) and [child relations](#) respectively.

anonymous function python also allows to make [anonymous functions](#) via the lambda constructor for [function objects](#):

```
lambda (p1, ..., pn): «expr»
```

anti-reflexive Defined along with [reflexive](#)

antisymmetric Defined along with [symmetric](#)

antisymmetric Defined along with [reflexive](#)

application Defined along with [partial function](#)

application layer The [application layer](#) of the internet protocol suite contains all protocols and methods that fall into the realm of process-to-process communications via an Internet Protocol (IP) network using the Transport Layer protocols to establish underlying host-to-host connections.

programming interface An [application programming interface](#) (API) is an [interface](#) that defines interactions between multiple software intermediaries. It defines the kinds of calls or requests that can be made, how to make them, the data formats that should be used, the conventions to follow, etc.

arbitrary Defined along with [variable](#)

arc Defined along with [graph](#)

architectural work Defined along with [copyrightable work](#)


argument Defined along with [subroutine](#)

argument Defined along with [partial function](#)

arithmetic/logic unit Defined along with [central processing unit](#)

arity A python **function** is defined by a code snippet of the form

```
def f (p1, ..., pn):  
    """docstring, what does this function do on parameters  
       :param pi: document arguments}  
    """  
    <<body>> # it can contain p1, ..., pn, and even f  
    return <<value>> # value of the function call (e.g text or number)  
<<more code>>
```

- the indented part is called the **body** of f , (: whitespace matters in python)
- the p_i are called **parameters**, and n the **arity** of f .

A function f can be **called** on **arguments** a_1, \dots, a_n by writing the expression $f(a_1, \dots, a_n)$. This executes the body of f where the (formal) parameters p_i are replaced by the arguments a_i .

arity Defined along with [subroutine](#)

assign Defined along with [variable assignment](#)

associative array See [dictionary](#)

asymmetric Defined along with [symmetric](#)

asymmetric Defined along with [reflexive](#)

asymptotically bounded Let $f, g: \mathbb{N} \rightarrow \mathbb{N}$, we say that f is **asymptotically bounded** by g , written as $f \leq_a g$, iff there is an $n_0 \in \mathbb{N}$, such that $f(n) \leq g(n)$ for all $n > n_0$.

asymptotically bounded Let $f, g: \mathbb{N} \rightarrow \mathbb{N}$, we say that f is **asymptotically bounded** by g (write $f \leq_a g$), iff there is an $n_0 \in \mathbb{N}$, such that $f(n) \leq g(n)$ for all $n > n_0$.

atto Defined along with [prefixes](#)

attribute Defined along with [opening tag](#)

attribute Defined along with [object-oriented programming](#)

attribute node Defined along with [XML document tree](#)

audiovisual work Defined along with [copyrightable work](#)

aural markup Defined along with [visual markup](#)

authority Defined along with [uniform resource identifier](#)

auxiliary storage See [secondary storage](#)

axiom An **axiom** (or **postulate**) is a statement about mathematical objects that we **assume to be true**.

base Defined along with [positional number system](#)

base Defined along with [positional number system](#)

base Defined along with [exponentiation](#)

base Defined along with [positional number system](#)

base Defined along with [exponentiation](#)

base Defined along with [exponentiation](#)

base Defined along with [exponentiation](#)

base equation **Error: The def does not appear to be inside a definition environment. line 75:3**

base name Defined along with [file system](#)

base set We call a structure $\langle S, \leq \rangle$ of a set S (the **base set**) equipped with a [preorder](#) r an **preordered set** or **proset**.

base unit Defined along with [International System of Units](#)

basic multilingual plane Most (non-Chinese) characters have code points in $[1, 65536]$ (the **basic multilingual plane**).

begin tag Defined along with [tag](#)

bijection Defined along with [bijjective](#)

bijjective Defined along with [injective](#)

bijjective A [function](#) $f: S \rightarrow T$ is called **bijjective** (or a **bijection** or a **one-to-one correspondence**), iff f is [injective](#) and [surjective](#).

binary Defined along with [source](#)

binary Defined along with [unary](#)

binary A code is called **binary** iff $B = \{0, 1\}$.

binary file Defined along with [text file](#)

binary unit prefix The following **binary unit prefix** es are used for units of information because they are similar to the [SI unit prefixes](#).

prefix	symbol	2^n	decimal	~SI prefix	Symbol
kibi	Ki	2^{10}	1024	kilo	k
mebi	Mi	2^{20}	1048576	mega	M
gibi	Gi	2^{30}	1.074×10^9	giga	G
tebi	Ti	2^{40}	1.1×10^{12}	tera	T
pebi	Pi	2^{50}	1.125×10^{15}	peta	P
exbi	Ei	2^{60}	1.153×10^{18}	exa	E
zebi	Zi	2^{70}	1.181×10^{21}	zetta	Z
yobi	Yi	2^{80}	1.209×10^{24}	yotta	Y

bit A **bit** (a contraction of “binary digit”) is the basic [unit](#) of capacity of a data storage device or communication channel. The capacity of a system which can exist in only two states, is one **bit** (written as 1 b)

blank symbol Defined along with [nondeterministic Turing machine](#)

body Defined along with [loop](#)

body Defined along with [function definition](#)

body Defined along with [subroutine](#)

bound An occurrence of a **variable** v is called **bound** in an **expression** E , iff it is in a **variable binding** that binds v . **Variables** that are not **bound** in an **expression** E are called **free** in E . We often write an **expression** E in which **variables** x_1, \dots, x_n occur **freely**, as $E[x_1, \dots, x_n]$.

branch A **path** in a **tree** that **starts** with the **root** is called a **branch**.

branch Defined along with **graph**

branches Defined along with **conditional execution**

branching factor Defined along with **in-degree**

browsing Defined along with **hypertext**

byte The **byte** is a derived **unit** for information capacity: $1 \text{ B} = 8 \text{ b}$.

calendar A **calendar** is a system of organizing **dates**. This is done by giving names to periods of time, typically **days**, **weeks**, **month** and **years**. A **date** is the designation of a single, specific **day** within such a system.

called Defined along with **subroutine**

canonical projection Defined along with **equivalence class**

canonical surjection Defined along with **equivalence class**

cardinality Defined along with **finite**

cell **Jupyter notebooks** consist of **cells** which come in three forms

- a **raw cell** shows text as is
- a **markdown cell** interprets the contents as markdown text (later more)
- a **code cell** interprets the contents as (e.g. **python**) code

cell Defined along with **Turing machine**

centi Defined along with **prefixes**

central processing unit A **central processing unit** (**CPU**), also called a **central processor** or **main processor**, is the electronic circuitry within a **computer** that carries out the **instructions** of a **program** by performing the basic arithmetic, logic, controlling, and input/output (I/O) operations specified by the **instructions**.

A **CPU** that consists of a

- **control unit** that interprets the **program** and controls the flow of **machine instructions** and
- a **arithmetic/logic unit** (**ALU**) that does the actual computations internally.

central processor See **central processing unit**

character Defined along with **universal character set**

character Defined along with **alphabet**

character code Let A and B be alphabets, then we call an injective function $c: A \rightarrow B^+$ a **character code**. A string $c(w) \in \{c(a) \mid a \in A\}$ is called a **codeword**.

character encoding A **character encoding** is a mapping from bit strings to UCS code points.

- character name** Defined along with [code point](#)
- character properties** Defined along with [code point](#)
- characteristic** Defined along with [floating point number](#)
- child** Defined along with [tree](#)
- child** Defined along with [tree](#)
- choreographic work** Defined along with [copyrightable work](#)
- civil law tradition** Legal systems in the [civil law tradition](#) are usually based on explicitly codified laws (civil codes).
- class** In [object-oriented programming](#), a [class](#) is a program construct for creating [objects](#) as well as providing the [fields](#) with initial [values](#) and the [methods](#) with implementations.
- client** A [client](#) is a piece of [computer hardware](#) or [software](#) that accesses a [service](#) made available by a [server](#).
- client-server architecture** In the [client-server architecture](#) a single overall computation is distributed across multiple [processes](#) or [computers](#).
A single [server](#) can serve multiple [clients](#), and a single [client](#) can use multiple [servers](#). A [client](#) may run on the same [computer](#) or may connect over a network to a [server](#) on a different [computer](#).
- closed** Defined along with [opened](#)
- closed** An [expression](#) is called [closed](#) or [ground](#), iff it does not contain [free variables](#).
- closing tag** Defined along with [opening tag](#)
- cloud IDE** See [web IDE](#)
- code** See [computer code](#)
- code** A [code](#) is a system of rules to convert [information](#) – such as a letter, word, sound, image, or gesture – into another form, sometimes shortened or secret, for communication via a [communication medium](#) or storage in a [storage medium](#).
The process of [encoding](#) applies a [code](#) for communication or storage, whereas the process of [decoding](#) applies it in reverse to restore the original [information](#)
- code cell** Defined along with [cell](#)
- code point** Each [UCS character](#) is identified by an unambiguous name and an integer number called its [code point](#).
- code point** For each character UniCode defines a [code point](#) (a number written in [hexadecimal](#) as U+ *ABCD*), a [character name](#), and a set of [character properties](#).
- codeword** Defined along with [character code](#)
- coding** The implementation of an [algorithm](#) in a chosen [programming language](#) is called [coding](#).
- codomain** Defined along with [domain](#)
- codomain** Defined along with [partial function](#)

- coefficient** Defined along with [floating point number](#)
- coefficient** Defined along with [scientific notation](#)
- collection** See [container](#)
- command** Defined along with [command-line interface](#)
- command line** Defined along with [command-line interface](#)
- command-line interface** A [command-line interface](#) (CLI) is a means of interacting with a [computer program](#) where the [user](#) (or [client](#)) issues [instructions](#) (called [commands](#) in a CLI) to the [program](#) in the form of successive lines of text ([command line](#)). The [program](#) which handles this [user interface](#) is called a [command-line interpreter](#) or [command-line processor](#).
- command-line interpreter** Defined along with [command-line interface](#)
- command-line processor** Defined along with [command-line interface](#)
- common law tradition** Legal systems in the [common law tradition](#) are usually based on case law, they are often derived from the British system.
- mathematical language** Defined along with [formulae](#)
- communication** [Communication](#) is the act of conveying [information](#) from one group of subjects to another.
- communication medium** A [communication medium](#) is a channel or system of [communication](#) – the means by which [information](#) (the [message](#)) is transmitted between a speaker or writer (the [sender](#)) and an audience (the [receiver](#)).
- compile** Defined along with [compiler](#)
- compiler** A [compiler](#) is a [program](#) that translates ([compiles](#))[code](#) written in one [programming language](#) (the [source language](#)) into another [language](#) (the [target language](#)).
- complement** See [set difference](#)
- complex numbers** The set \mathbb{C} of [complex numbers](#) contains [expressions](#) of the form $c = (a + bi)$, where $a, b \in \mathbb{R}$. $\text{Re}(c) := a$ is called the [real part](#) and $\text{Im}(c) := b$ the [imaginary part](#) of c ; we call i the [imaginary unit](#).
- complexes** Defined along with [integer](#)
- component** Defined along with [uniform resource identifier](#)
- component** Defined along with [mathematical structure](#)
- composable** We call two relations $R \subseteq A \times B$ and $S \subseteq C \times D$ [composable](#), iff $B = C$
- compose** Defined along with [composition principle](#)
- composition** The [composition](#) of two relations $R \subseteq A \times B$ and $S \subseteq B \times C$ is defined as

$$S \circ R := \{(a, c) \in A \times C \mid \text{there is a } b \in B \text{ with } (a, b) \in R \text{ and } (b, c) \in S\}$$
- composition** The [composition](#) of $R \subseteq A \times B$ and $S \subseteq B \times C$ is defined as $S \circ R := \{(a, c) \in A \times C \mid \exists b \in B. (a, b) \in R \wedge (b, c) \in S\}$
- composition principle** All [programming languages](#) provide [composition principles](#) that allow to [compose](#) smaller program fragments into larger ones in such a way, that the [semantics](#) of the larger is determined by the [semantics](#) of the smaller ones and that of the [composition principle](#) employed.

- putationally universal** See [Turing complete](#)
- computer** See [computing device](#)
- computer code** We call any well-formed fragments of a [program computer code](#) or [program code](#), or just [code](#).
- computer programming** [Computer programming](#) (or just [programming](#)) is the process of designing and building a [program](#) for accomplishing a specific computing task.
It involves sub-processes, such as: analysis, generating [algorithms](#), profiling [algorithms'](#) resource consumption, proving [algorithm](#) properties, [coding](#), and program verification.
- computer science** [Computer science](#) (or short [CS](#)) is the study of [algorithms](#) and [information processing system](#) in theory and practice. A [CS](#) professional is called a [computer scientist](#).
- computer scientist** Defined along with [computer science](#)
- computing device** A [computing device](#) or simply a [computer](#) is an physical (usually electrical or electronic) [information processing system](#) that can automatically [execute](#) a [sequence](#) of [machine instructions](#) i.e. arithmetic or logical operations that change [state](#) of the [system](#).
A [computer](#) consists of physical parts (its [hardware](#)) and a set of [programs](#) and [data](#), its [software](#).
- concatenation** The [concatenation](#) $\text{conc}(L_1, L_2)$ of two languages L_1 and L_2 over the same alphabet is defined as $\text{conc}(L_1, L_2) := \{s_1s_2 \mid s_1 \in L_1 \wedge s_2 \in L_2\}$.
- concatenation** The [concatenation](#) $\text{conc}(s, t)$ of two strings $s = \langle s_1, \dots, s_n \rangle \in A^n$ and $t = \langle t_1, \dots, t_m \rangle \in A^m$ is defined as $\langle s_1, \dots, s_n, t_1, \dots, t_m \rangle \in A^{(n+m)}$.
We will often write $\text{conc}(s, t)$ as $s + t$ or simply st .
- condition** A [condition](#) is a [Boolean expression](#) in a [control structure](#).
- conditional execution** [Conditional execution](#) allows to execute (or not to execute) certain parts of a [program](#) (the [branches](#)) depending on a [condition](#). We call a code block that enables [conditional execution](#) a [conditional statement](#).
- conditional statement** Defined along with [conditional execution](#)
- conjecture** **Error: The def does not appear to be inside a definition environment. line 43:51**
- connective** Defined along with [formulae](#)
- constant** A [constant](#) is a [memory](#) location which contains a [value](#) that cannot be altered by the program during normal execution. It is [referenced](#) by an [identifier](#) – the [constant name](#).
- constant name** Defined along with [constant](#)
- container** A [container](#) or [collection](#) is a grouping of some variable number (possibly zero) of [data](#) items – the [elements](#) of the [container](#) – that need to be operated upon together in some controlled fashion.
- control flow** The [control flow](#) of a [program](#) is the sequence of execution of the [program instructions](#). It is specified via special [program instructions](#) called [control structures](#).
- control structure** Defined along with [control flow](#)

control unit Defined along with [central processing unit](#)

control word Defined along with [document markup](#)

converse $R^{-1} := \{(y, x) \mid (x, y) \in R\}$ is the **converse** relation of R .

converse relation $R^{-1} := \{(y, x) \mid (x, y) \in R\}$ is the **converse relation** of $R \subseteq A \times B$.

copyright Defined along with [intellectual property](#)

copyright holder The **copyright holder** is the legal entity that owns the [copyright](#) to a [copyrighted](#) work.

copyright infringement The use of a [copyrighted](#) material, by anyone other than the owner of the [copyright](#), amounts to **copyright infringement** only when the use is such that it conflicts with any one or more of the exclusive rights conferred to the owner of the [copyright](#).

copyrightable work A **copyrightable work** is any artefact of human labor that fits into one of the following eight categories:

- **Literary works**: Any work expressed in letters, numbers, or symbols, regardless of medium. (Computer source code is also considered to be a literary work.)
- **Musical works**: Original musical compositions.
- **Sound recordings** of musical works. (different licensing)
- **Dramatic works**: literary works that direct a performance through written instructions.
- **Choreographic works** must be “fixed,” either through notation or video recording.
- **Pictorial, graphic and sculptural work (PGS works)**: Any two-dimensional or three-dimensional art work
- **Audiovisual works**: work that combines audio and visual components. (e.g. films, television programs)
- **Architectural works**. (copyright only extends to aesthetics)

copyrighted Defined along with [public domain](#)

corollary **Error: The def**
defi does not appear to be inside a definition environment. line 49:3

countable We say that a [set](#) A is **countable** (otherwise **uncountable**), iff there is an [bijjective](#) function $f: A \rightarrow \mathbb{N}$ with $N \subseteq \mathbb{N}$.

countably infinite We say that a [set](#) A is **countably infinite**, iff there is a [bijjective](#) function $f: A \rightarrow \mathbb{N}$.

current Defined along with [length](#)

cycle Defined along with [cyclic](#)

cycle Defined along with [cyclic](#)

cyclic Given a graph $G = \langle V, E \rangle$, then

- a [path](#) p is called **cyclic** (or a **cycle**) iff $\text{start}(p) = \text{end}(p)$.
- a cycle $\langle v_0, \dots, v_n \rangle$ is called **simple**, iff $v_i \neq v_j$ for $1 \leq i, j \leq n$ with $i \neq j$.
- graph G is called **acyclic** iff there is no cyclic path in G .

cyclic Given a [directed graph](#) $G := \langle V, E \rangle$,

- a path p is called **cyclic** (or a **cycle**) iff $\text{start}(p) = \text{end}(p)$.
- a cycle $\langle v_0, \dots, v_n \rangle$ is called **simple**, iff $v_i \neq v_j$ for $1 \leq i, j \leq n$ with $i \neq j$.
- G is called **acyclic** (or a **DAG: directed acyclic graph**) iff there is no cycle in G .

dashboard A **dashboard** is a **user interface** that organizes and presents **system information** give an overview over the **state** of a complex **system** and its **services**.

data **Data** is **information** that is used to represent objects by giving values to their relevant attributes and stating their relationships.

data language A **data language** is a **formal language** for specifying **data** in an **information processing system**. **data languages** are not **Turing complete**.

data structure A **data structure** is a **data** organization, management, and storage format that enables efficient access and modification.

data type A **data type** or simply **type** is an **programming language** attribute of **data** which tells the **compiler** or **interpreter** how the **programmer** intends to use the **data**.

A **type** constrains the **values** a **variable** or **function** can might take and defines the **operators** that can be applied to it.

date Defined along with **calendar**

day Defined along with **minute**

deca Defined along with **prefixes**

deci Defined along with **prefixes**

decimal Defined along with **unary**

decimal point Defined along with **floating point number**

decoding Defined along with **code**

namespace declaration Defined along with **namespace declaration**

default value Defined along with **keyword argument**

defined at We call a **partial function** $f: X \rightarrow Y$

- **defined at** $x \in X$, iff $(x, y) \in f$ for some $y \in Y$ and
- **undefined at** $x \in X$ (write $f(x) = \perp$), iff $(x, y) \notin f$ for all $y \in Y$.

defined piecewise A **function** m is **defined piecewise**, we write

$$m(x) = \begin{cases} a_1 & \text{if } A_1 \\ \vdots & \vdots \\ a_n & \text{if } A_n \\ o & \text{else} \end{cases}$$

where A_i are conditions involving x , if $m(x) = a_i$ for all x with A_i and o otherwise.

definiendum **Error: The defi does not appear to be inside a definition environment. line 22:26**

definiendum Defined along with **definitional equation**

- definiens** **Error: The def does not appear to be inside a definition environment. line 23:3**
- definiens** Defined along with [definitional equation](#)
- definite integral** Given a function $f: \mathbb{R} \rightarrow \mathbb{R}$ and an interval $[a, b] \subset \mathbb{R}$, then the **definite integral** $\int_b^a f(x)dx$ is defined to be the signed area of the region in the plane bounded by the graph of f , the x -axis, and the vertical lines $x = a$ and $x = b$, such that area above the x -axis adds to the total, and that below the x -axis subtracts from the total.
- definition by description** **Error: The def does not appear to be inside a definition environment. line 58:16**
- definitional equation** If a does not occur in A , we call a pair $a := A$ a **definitional equation** and $a :\Leftrightarrow A$ a **definitional equivalence** with **definiendum** a and **definiens** A .
- definitional equivalence** Defined along with [definitional equation](#)
- denominator** Defined along with [rational number](#)
- depth** The **depth** of a [node](#) n in a [tree](#) t is the [length](#) of the [path](#) that [links](#) n with the [root](#) of t .
- dereferencing** Defined along with [reference](#)
- derivation** Defined along with [inference](#)
- derived unit** A **derived unit** is formed as a product of integer powers of [base units](#).
- descendant** Defined along with [tree](#)
- descendant** Defined along with [ancestor](#)
- nondeterministic Turing machine** Defined along with [nondeterministic Turing machine](#)
- diagonal** See [identity function](#)
- dictionary** A **dictionary** (also called [associative array](#), [map](#), [symbol table](#)) is an [abstract data type](#) composed of a [set](#) of [key/value](#) pairs, such that each possible [key](#) appears at most once in the [container](#).
- difference** Defined along with [subtraction](#)
- difference** Defined along with [subtraction](#)
- difference** Defined along with [subtraction](#)
- difference** Defined along with [subtraction](#)
- digit** Defined along with [positional number system](#)
- digit** Defined along with [positional number system](#)
- digit** Defined along with [positional number system](#)
- digital text** An [electronic document](#) that contains a digital encoding of textual material that can be read by the [end user](#) by simply presenting the encoded [characters](#) is called **digital text**.
- digraph** Defined along with [graph](#)
- digraph** Defined along with [directed graph](#)

dimension Defined along with [n-fold Cartesian product](#)

dimension Defined along with [International System of Units](#)

directed acyclic graph Defined along with [cyclic](#)

directed edge Defined along with [directed graph](#)

directed graph Defined along with [graph](#)

directed graph A [directed graph](#) (also called [digraph](#)) is a [pair](#) $\langle V, E \rangle$ such that

- V is a set of vertices
- $E \subseteq V \times V$ is the set of its [directed edges](#)

division The [division](#) operator computes the [quotient](#) a/b of $a \in \mathbb{Q}$ and $b \in \mathbb{Q}$. On \mathbb{Q} we define $\frac{a}{b} / \frac{c}{d} := \frac{a \cdot d}{b \cdot c}$.

division [Division](#) computes the [modulus](#) $n \operatorname{div} m$ of two [natural numbers](#) n and m . $n \operatorname{div} m$ is defined as that $q \in \mathbb{N}$, such that $n = m \cdot q + r$ for some $r \leq 0 < m$. The number r is called the [remainder](#) and is written as $n \operatorname{mod} m$.

division [Division](#) computes the [quotient](#) a/b of a and b . It is defined as is that c – if it exists, such that $a \cdot c = b$.

division Defined along with [absolute value](#)

document format A [document format](#) is a [file format](#) for [electronic documents](#).

document markup [Document markup](#) (or just [markup](#)) is the process of adding [control words](#) (special character sequences also called [markup codes](#)) to a [plain text](#) to control the structure, formatting, or the relationship among its parts, making it a [formatted text](#). All [characters](#) of a [formatted text](#) that are not [control words](#) constitute its [textual content](#).

document object model The [document object model](#) ([DOM](#)) is a [data structure](#) for storing [marked-up electronic documents](#) as [trees](#) together with a standardized set of access methods for manipulating them.

document renderer Defined along with [electronic document](#)

document root As a document is a [tree](#), the [XML](#) specification mandates that there must be a unique [document root](#).

document type Defined along with [markup format](#)

domain call X the [domain](#) (write $\operatorname{dom}(f)$), and Y the [codomain](#) ($\operatorname{codom}(f)$) (come with f)

domain Defined along with [partial function](#)

dot notation Defined along with [object](#)

double star operator The [double star operator](#) unpacks a [dictionary](#) into a sequence of [keyword arguments](#).

dramatic work Defined along with [copyrightable work](#)

edge Defined along with [graph](#)

electronic document An [electronic document](#) is any electronic [media content](#) that is intended to be used via a [document renderer](#), i.e. a [program](#) or [computing device](#) that transforms it into a form that can be directly perceived by the [end user](#).

element Defined along with [set](#)

element Defined along with [set](#)

element Defined along with [container](#)

element node Defined along with [XML document tree](#)

elementhood Defined along with [set](#)

embedded system An **embedded system** is a [computing device](#) with a dedicated function within a larger mechanical or electrical [system](#).

empty Defined along with [empty set](#)

empty element tag Defined along with [opening tag](#)

empty set the **empty set**: $\forall x. x \notin \emptyset$

empty set The **empty set** \emptyset (also written as $\{\}$) is the [set](#) without elements. A set is called **empty**, iff it is \emptyset , and **non-empty** or (**inhabited**) otherwise.

empty string Defined along with [alphabet](#)

empty tag Defined along with [tag](#)

empty word Defined along with [alphabet](#)

encoding Defined along with [code](#)

encoding scheme UniCode defines various **encoding schemes** for characters, the most important is UTF-8.

end Defined along with [path](#)

end Defined along with [path](#)

end tag Defined along with [tag](#)

end user An **end user** is a person who ultimately uses or is intended to ultimately use a [computation device](#) or [program](#).

end user license agreement Software vendors usually license software under extensive **end-user license agreement** (EULA) entered into upon the installation of that software on a computer. The license authorizes the user to install the software on a limited number of computers.

equal Two sets A and B are **equal** (written $A \equiv B$), iff they have the same elements.

equal We call two mathematical objects a and b **equal**, (written $a = b$), iff there are no properties that discern them.

equivalence class Let S be a set and R be an equivalence relation on S , then for any we call $x \in S$ we call the set $[x]_R := \{y \in S \mid R(x, y)\}$ the **equivalence class** of x (under R), and the set $S/R := \{[x]_R \mid x \in S\}$ the **quotient space** of S (under R).

equivalence class Let S be a set and R be an [equivalence relation](#) on S , then for any we call $x \in S$ we call the set $[x]_R := \{y \in S \mid R(x, y)\}$ the **equivalence class** of x (under R), and the set $S/R := \{[x]_R \mid x \in S\}$ the **quotient space** or **quotient set** of S (under R), it is often read as S “**modulo** R ”. The element x is called the **representative** of $[x]_R \in S/R$.

The mapping $\pi_R: S \rightarrow S/R; x \mapsto [x]_R$ is called the **canonical projection** or **canonical surjection** of S to S/R .

equivalence relation Defined along with [reflexive](#)

equivalence relation A relation $R \subseteq A \times A$ is an [equivalence relation](#) on A , iff R is [reflexive](#), [symmetric](#), and [transitive](#).

escape character Defined along with [string literal](#)

escape sequence Defined along with [string literal](#)

evaluation Defined along with [expression](#)

exa Defined along with [prefixes](#)

exbi Defined along with [binary unit prefix](#)

execute Defined along with [interpreter](#)

execute Defined along with [computing device](#)

exploitation rights Defined along with [personal rights](#)

exponent Defined along with [floating point number](#)

exponent Defined along with [exponentiation](#)

exponent Defined along with [exponentiation](#)

exponent Defined along with [exponentiation](#)

exponent Defined along with [exponentiation](#)

exponentiation The [exponentiation](#) operation raises a number a (the [base](#)) to the [power](#) n (the [exponent](#)). We define it as

$$a^n := \begin{cases} b^n & \text{if } a \in \mathbb{Z}^- \text{ and } (n = 2 \cdot k) \\ -(b^n) & \text{if } a \in \mathbb{Z}^- \text{ and } (n = 2 \cdot k + 1) \end{cases}$$

exponentiation [Exponentiation](#) raises a number $a \in \mathbb{Q}$ (the [base](#)) to the [power](#) $n \in \mathbb{Q}$ (the [exponent](#)). We define

$$\frac{a}{b}^{\frac{n}{m}} := \frac{\sqrt[m]{a^n}}{\sqrt[m]{b^n}}$$

exponentiation [Exponentiation](#) raises a [natural number](#) a (the [base](#)) to the n -th [power](#) a^n ($n \in \mathbb{N}$ is called the [exponent](#)). We define $a^0 := 1$ and $a^{s(n)} := aa^n$.

exponentiation [Exponentiation](#) raises a [number](#) a (the [base](#)) to the n -th [power](#) a^n (n is called the [exponent](#)).

exponentiation Defined along with [absolute value](#)

expression An [expression](#) in a [programming language](#) is a combination of one or more [constants](#), [variables](#), [operators](#), and [functions](#) that the [programming language](#) computes to produce a [value](#). This process is called [evaluation](#).

expression An [expression](#) is a finite construction composed of [variables](#) and names of mathematical objects/concepts composed by [operator application](#) and [variable binding](#) according to rules that depend of the (mathematical) context.

extension The [extension](#) of a code (on characters) $c: A \rightarrow B^+$ to a function $c': A^* \rightarrow B^*$ is defined as $c'(\langle a_1, \dots, a_n \rangle) = \langle c(a_1), \dots, c(a_n) \rangle$.

extension Defined along with [file system](#)

external memory See [secondary storage](#)

f-string See [formatted string literal](#)

f-string See [formatted string literal](#)

fair use doctrine Case law in [common law traditions](#) has established a [fair use doctrine](#), which allows e.g.

- making safety copies of software and audiovisual data,
- lending of books in public libraries,
- citing for scientific and educational purposes, or
- excerpts in search engine.

Fair use is established in court on a case-by-case taking into account the purpose (commercial/educational), the nature of the work the amount of the excerpt, the effect on the marketability of the work.

false Defined along with [truth value](#)

falsum Defined along with [truth value](#)

femto Defined along with [prefixes](#)

field Defined along with [object-oriented programming](#)

file A [file](#) is a [resource](#) for recording [data](#) in a [storage device](#).

file format A [file format](#) is a standard way that [information](#) is [encoded](#) for storage in a computer [file](#). It specifies how bits are used to encode information in a [storage device](#).

file object python uses [file objects](#) to encapsulate all file input/output functionality.

file system A [file system](#) is a [program](#) that organizes space on a [storage device](#) and makes it accessible as [files](#). A [file name](#) usually consists of a [base name](#) and an [extension](#) separated by a dot.

final state Defined along with [nondeterministic Turing machine](#)

finite We say that a set A is [finite](#) and has [cardinality](#) (or [size](#)) $\#(A) \in \mathbb{N}$, iff there is a [bijective](#) function $f: A \rightarrow \{n \in \mathbb{N} \mid n < \#(A)\}$.

The cardinality of a set A is also written as $|A|$, $\text{card}(A)$, $n(A)$, or \overline{A} .

finite sequence Defined along with [sequence](#)

first component Let $p := (a, b)$ be a pair, then we call $p^1 := a$ the [first component](#) and $p^2 := b$ the [second component](#) of p .

float Defined along with [integer](#)

float See [floating point number](#)

floating point number A [floating point number](#) (or short a [float](#)) is a quintuple $n := \langle \sigma, s_1, s_2, b, e \rangle$, where

1. the [sign](#) σ is unit sequence – or the empty sequence, and the s_i are [sequence](#) of [digits](#) of [base](#) b . Together, (i.e. concatenated with a [decimal point](#) between s_1 and s_2) σ, s_1 , and s_2 make up the [significand](#) s (also called the [mantissa](#) or [coefficient](#)).
2. an [exponent](#) $e \in \mathbb{Z}$ (also referred to as the [characteristic](#), or [scale](#)), which modifies the magnitude of the number n .

The number $\langle \sigma, s_1, s_2, b, e \rangle$ represents the rational number $\frac{s}{b^{p-1}} \cdot b^e$.

The **length** $p := \text{len}(s_1) + \text{len}(s_2)$ of the **significand** determines the **precision** to which numbers can be represented.

for loop A **for loop iterates** a **program** fragment over a **sequence**; we call the process **iteration**. python uses the following general syntax

```
for <<var>> in <<range>>:  
    <<body>>  
<<other code>>
```

form action Defined along with **input element**

form data The **HTML** form element groups the layout and input elements:

- `<form action="<<URI>>"method="<<req>>">` specifies the **form action** in terms of a **HTTP request** $\langle \text{req} \rangle$ to the **URI** $\langle \text{URI} \rangle$.
- The **form data** consists of a string $\langle \text{data} \rangle$ of the form $n_1=v_1 \& \dots \& n_k=v_k$, where
 - * n_i are the values of the **name** attributes of the input fields
 - * and v_i are their values at the time of submission.
- `<input type="submit" .../>` triggers the **form action**: it composes a **HTTP request**
 - * If $\langle \text{req} \rangle$ is **get** (the default), then the browser issues a **GET** request $\langle \text{URI} \rangle? \langle \text{data} \rangle$.
 - * If $\langle \text{req} \rangle$ is **post**, then the browser issues a **POST** request to $\langle \text{URI} \rangle$ with document content $\langle \text{data} \rangle$.

formal language A set $L \subseteq A^*$ is called a **formal language** in A .

formal language A **formal language** (or simply **language**) over an **alphabet** \mathcal{A} is a set $\mathcal{L} \subseteq \mathcal{A}^*$ of **words** over \mathcal{A} .

formatted string literal **Formatted string literals** (aka. **f-strings**) are **string literals** can contain python expressions that will be replaced with their values at runtime.

F-strings are prefixed by a prefix **f** or **F**, the expressions are delimited by curly braces, and the characters **{** and **}** themselves are represented by **{{** and **}}**.

formatted text Defined along with **plain text**

formulae Mathematicians use a stylized language that

- uses **formulae** to represent mathematical objects, e.g. $\int_1^0 x^{3/2} dx$
- uses **math idioms** for special situations (e.g. *iff, hence, let...be..., then...*)
- classifies statements by role (e.g. **Definition, Lemma, Theorem, Proof, Example**)

We call this language **mathematical vernacular** or **common mathematical language**.

For the use in formulae we use abbreviations (special symbols) for many of the **connectives**:

- \wedge and \vee and \neg are common notations for **and** and **or**
- **not** is in mathematical statements often denoted with \neg
- **iff** as abbreviation for **if and only if**, symbolized by \Leftrightarrow
- the symbol \Rightarrow is used as a shortcut for **implies** or **if..., then...**

and **quantifiers**:

- $\forall x.P$ ($\forall x \in S.P$) stands for *P holds for all x (in S)*
- $\exists x.P$ ($\exists x \in S.P$) stands for *there exists an x (in S) such that P holds*
- $\nexists x.P$ ($\nexists x \in S.P$) stands for *there exists no x (in S) such that P holds*
- $\exists^1 x.P$ ($\exists^1 x \in S.P$) stands for *there exists one and only one x (in S) such that P holds*

fraction Defined along with [rational number](#)

fragment identifier Defined along with [uniform resource identifier](#)

free Defined along with [bound](#)

function Defined along with [subroutine](#)

function If we do not want to specify whether a partial function is [total](#), then we simply speak of a [function](#).

function definition The [function definition](#) $f(a_1, \dots, a_n) := B[a_1, \dots, a_n]$ defines a n -ary [function](#) f by its behavior on the (formal) [arguments](#) a_1, \dots, a_n . Here we call $f(a_1, \dots, a_n)$ the [function pattern](#) and $B[a_1, \dots, a_n]$ the [body](#)—an [expression](#) that can contain the [arguments](#) a_1, \dots, a_n as [free variables](#).

A [relation](#) p is defined analogously via [definitional equivalence](#) $p(a_1, \dots, a_n) :\Leftrightarrow B[a_1, \dots, a_n]$.

function object Defined along with [anonymous function](#)

function pattern Defined along with [function definition](#)

function space Given sets A and B We will call the set $A \rightarrow B$ ($A \multimap B$) of all (partial) functions from A to B the (partial) [function space](#) from A to B .

function space Defined along with [partial function space](#)

general-purpose computer A [general-purpose computer](#) is [one](#) that, given the appropriate [Software](#) and the required time, should be able to perform arbitrary computing tasks.

gibi Defined along with [binary unit prefix](#)

giga Defined along with [prefixes](#)

graph A [graph](#) is a [pair](#) $G := \langle V, E \rangle$ such that V is a set and $E \subseteq V \times V$ is a [relation on V](#). We call V the [vertices](#) (or [nodes](#), [points](#), [junctions](#)) and E the [edges](#) (or [lines](#), [branches](#), [arcs](#)) of G .

If E is [symmetric](#), we call G an [undirected graph](#), else a [directed graph](#) or [digraph](#). In the former we consider the pairs $(a, b), (b, a) \in E$ together as an [unordered pair](#) $\{a, b\}$.

graphical user interface A [graphical user interface](#) is a [user interface](#) that includes graphical elements, such as windows, icons, and buttons.

greater than Defined along with [less than](#)

ground See [closed](#)

halt Defined along with [Turing machine](#)

handle See [identifier](#)

hardware Defined along with [computing device](#)

head Defined along with [Turing machine](#)

hecto Defined along with [prefixes](#)

height The **height** of a **node** n in a **tree** t is the **length** of the longest **path** that **links** n to a **leaf** of t . The **height** of t is the **height** of its **root**.

height Defined along with [height](#)

hexadecimal Defined along with [unary](#)

higher-order function We call a **function** a **higher-order function**, iff it takes a **function** as **argument**.

hour Defined along with [minute](#)

hyperlink A **hyperlink** is a reference to data that can immediately be followed by the user or that is followed automatically by a user agent.

hypertext A collection text documents with hyperlinks that point to text fragments within the collection is called a **hypertext**. The action of following hyperlinks in a hypertext is called **browsing** or **navigating** the hypertext.

idempotent We call a **HTTP** request **idempotent**, iff executing it twice has the same effect as executing it once.

identifier An **identifier** (also called **handle**) is **reference** to a **resource**.

identity function The **identity function** on a set A is defined as $\text{Id}_A := \{(a, a) \mid a \in A\}$.

identity function For a **set** A , the **identity function** $\text{Id}_A: A \rightarrow A$ on A maps any $a \in A$ to itself. If we think of Id_A as a **relation** on A , then we call it the **identity relation** or **diagonal** on A and write it as Δ_A .

identity relation See [identity function](#)

iff Defined along with [MathTalk](#)

iff Defined along with [formulae](#)

image Let $f: A \rightarrow B$ be a function, $A' \subseteq A$, and $B' \subseteq B$, then we call

- $f(A') := \{b \in B \mid (a, b) \in f \text{ for some } a \in A'\}$ the **image** of A' under f ,
- $\mathbf{Im}(f) := f(A)$ the **image** of f , and
- $f^{-1}(B') := \{a \in A \mid (a, b) \in f \text{ for some } b \in B'\}$ the **pre-image** of B' under f .

image Defined along with [image](#)

imaginary part Defined along with [complex numbers](#)

imaginary unit Defined along with [complex numbers](#)

immutable Defined along with [mutable](#)

implicit definition **Error: The def does not appear to be inside a definition environment. line 57:56**

in-degree Given a graph $G = \langle V, E \rangle$. The **in-degree** $\text{indeg}(v)$ and the **out-degree** $\text{outdeg}(v)$ (or **branching factor**) of a **vertex** $v \in V$ are defined as

- $\text{indeg}(v) = \#\{w \mid (w, v) \in E\}$
- $\text{outdeg}(v) = \#\{w \mid (v, w) \in E\}$

indegree Let $G := \langle V, E \rangle$ be a **directed graph** and $v \in V$ a **node** in G , then we define

- **indegree** $\text{indeg}(v)$ of v as $\#\{w \mid (w, v) \in E\}$
- **outdegree** $\text{outdeg}(v)$ of v as $\#\{w \mid (v, w) \in E\}$

induced undirected graph Let $G := \langle V, E \rangle$ be a **graph**, and E' the **symmetric closure** of E , then we call $\langle V, E' \rangle$ the **induced undirected graph** of G .

industrial design right Defined along with **intellectual property**

inference We call a sequence of **inferences** a **derivation** or a **proof** (of the last statement).

infinite A **set** that is not **finite** is called **infinite**.

infinite sequence Defined along with **sequence**

infinity **Infinity** (written ∞) is an abstract concept describing something without any limit. In mathematics is usually treated like a number.

information **Information** consists of a **sequence** of **symbols** or **states**.

information processing system An **information processing system** (or **information processor**) is a **stateful system** (be it electrical, mechanical or biological) which takes **information** in one form and transforms it into another form.

An **information processing system** S is made up of four **subsystems**:

1. the **input subsystem** channels **information** into S ,
2. the **processor** executes the transformation in a sequence of operations called **instructions** on the **processor state**,
3. the **storage subsystem** stores **information**, and
4. the **output subsystem** channels the transformed **information** out of S .

information processor See **information processing system**

inhabited Defined along with **empty set**

initial Let $G := \langle V, E \rangle$ be a **directed graph**, then we call a **node** $v \in V$

- **initial** (or a **source**) in G , iff there is no $w \in V$ such that $(w, v) \in E$.
- **terminal** (or a **sink**) in G , iff there is no $w \in V$ such that $(v, w) \in E$.

initial Let $G = \langle V, E \rangle$ be a **directed graph**, then we call a **node** $v \in V$

- **initial**, iff there is no $w \in V$ such that $(w, v) \in E$. (no predecessor)
- **terminal**, iff there is no $w \in V$ such that $(v, w) \in E$. (no successor)

In a graph G , node v is also called a **source** (**sink**) of G , iff it is initial (terminal) in G .

initial state Defined along with **nondeterministic Turing machine**

injection Given a **tuple** $v \in A_1 \times \dots \times A_{(i-1)} \times A_{(i+1)} \times \dots \times A_n$ we call the function $\iota_v^i: A_i \rightarrow A_1 \times \dots \times A_n; a_i \mapsto \langle a_1, \dots, a_n \rangle$ the (i^{th}) **injection** induced by v .

- injective** A function $f: S \rightarrow T$ is called
- **injective** iff $\forall x, y \in S. f(x) = f(y) \Rightarrow x = y$.
 - **surjective** iff $\forall y \in T. \exists x \in S. f(x) = y$.
 - **bijective** iff f is injective and surjective.

injective A **function** $f: S \rightarrow T$ is called **injective** or **one-to-one**, iff $f(x) = f(y)$ entails $x = y$ for all $x, y \in S$.

inner node Defined along with **path**

- input element** The **HTML** form element groups the layout and **input elements**:
- `<form action="⟨URI⟩"…>` specifies the **form action** (as a web page address).
 - `<input type="submit"…/>` triggers the **form action**: it sends the **form data** to web page specified there.

input subsystem Defined along with **information processing system**

input symbol Defined along with **nondeterministic Turing machine**

instruction Defined along with **information processing system**

integer python has the following five basic **data types**

Data type	Keyword	contains	Examples
integers	int	bounded integers	1, -5, 0, ...
floats	float	floating point numbers	1.2, .125, -1.0, ...
strings	str	strings	"Hello", 'Hello', "123", 'a', ...
Booleans	bool	truth values	True, False
complexess	complex	complex numbers	2+3j, ...

integer See **integer number**

integer division **Integer division** computes the **integer quotient** (or **modulus**) $n \text{ div } m$ of two **integers**. $n \text{ div } m$ is defined as that $q \in \mathbb{Z}$, such that $nm \cdot q + r$ for some $0 \leq r < m$. The number r is called the **remainder** and is written as $n \text{ mod } m$.

integer interval We define the **integer interval** as a set of consecutive integers: $[a, b] := \{x \in \mathbb{Z} \mid x \leq a \leq b\}$

integer number The set \mathbb{Z} of **integer numbers** (or **integers**) is defined as $\mathbb{Z} := \mathbb{N} \cup \{-n \mid n \in \mathbb{N}^+\}$.

integer quotient Defined along with **integer division**

development environment An **integrated development environment (IDE)** is a **program** that provides comprehensive facilities to computer programmers for software development. An **IDE** normally consists of at least a source code editor, build automation tools, and a debugger.

intellectual property The concept of **intellectual property** motivates a set of laws that regulate **property rights** rights on intangible objects, in particular

- **Patents** grant exploitation rights on original ideas.
- **Copyrights** grant personal and exploitation rights on expressions of ideas.
- **Industrial design rights** protect the visual design of objects beyond their function.
- **Trademarks** protect the signs that identify a legal entity or its products to establish brand recognition.

interactive toplevel See [read-eval-print loop](#)

interface See [user interface](#)

interpreter An [interpreter](#) is a [program](#) that directly [executes instructions](#) written in a [programming language](#), without requiring them previously to have been [compiled](#) into a machine language program.

intersection [intersection](#): $A \cap B := \{x \mid x \in A \wedge x \in B\}$

intersection Let A and B be sets, then the [intersection](#) $A \cap B$ of A and B is $\{x \mid x \in A \text{ and } x \in B\}$.

intersection Let I be a [set](#) and S_i a family of sets indexed by I , then the [intersection](#) $\bigcap_{i \in I} S_i$ over I is $\{x \mid x \in S_i \text{ for all } i \in I\}$.

intersection over a collection [intersection over a collection](#): Let I be a [set](#) and S_i a family of sets indexed by I , then $\bigcap_{i \in I} S_i := \{x \mid \forall i \in I. x \in S_i\}$.

interval We define four kinds of [intervals](#) as [subsets](#) of the [real numbers](#):

– $[a, b] := \{x \in \mathbb{R} \mid (a \leq x) \text{ and } (x \leq b)\}$

– $[a, b) := \{x \in \mathbb{R} \mid (a \leq x) \text{ and } (x < b)\}$

– $(a, b] := \{x \in \mathbb{R} \mid (a < x) \text{ and } (x \leq b)\}$

– $(a, b) := \{x \in \mathbb{R} \mid (a < x) \text{ and } (x < b)\}$

intransitive Defined along with [transitive](#)

inverse function If f is [bijective](#), call the [converse relation](#) [inverse function](#), we (also) write it as f^{-1} .

inverse function If $f: A \rightarrow B$ is [injective](#), then the [converse relation](#) is a partial function $f^{-1}: B \rightarrow A$, we call it the [inverse function](#) of f . If f is [bijective total](#) function, then f^{-1} is a [total function](#).

invoke Defined along with [subroutine](#)

invoker Defined along with [subroutine](#)

irreflexive Defined along with [reflexive](#)

irreflexive Defined along with [reflexive](#)

is (of complexity) Defined along with [Landau set](#)

iterate Defined along with [for loop](#)

iteration Defined along with [for loop](#)

junction Defined along with [graph](#)

key Defined along with [dictionary](#)

keyword argument The last $k \leq n$ of n parameters of a [function](#) can be [keyword arguments](#) of the form $p_i = \langle\langle \text{val} \rangle\rangle_i$: If no argument a_i is given in the function call, the [default value](#) $\langle\langle \text{val} \rangle\rangle_i$ is taken.

keyword argument python [functions](#) can take [keyword arguments](#): if k is a sequence of key/value pairs then `def f($p_1, \dots, p_n, **k$)`, binds the keys to values in the body of f .

kibi Defined along with [binary unit prefix](#)

kilo Defined along with [prefixes](#)

language See [formal language](#)

language shell See [read-eval-print loop](#)

leaf Defined along with [tree](#)

leaf Defined along with [tree](#)

leap year Defined along with [year](#)

lemma **Error: The def does not appear to be inside a definition environment. line 47:67**

length Defined along with [sequence](#)

length Defined along with [path](#)

length The [length](#) $|s|$ of a [word](#) $s \in A^n$ is n .

length Defined along with [path](#)

length The seven [dimensions](#) of [SI base units](#) are [length](#) (L), [mass](#) (M), [time](#) (T), [current](#) (I), [temperature](#) (Θ), [luminous intensity](#) (J), and [amount](#) (N).

less than We define the order relation $<_{\mathbb{Z}}$ (n is [less than](#) m also written as $<$) by

$$<_{\mathbb{Z}} := <_{\mathbb{N}} \cup \{(n, m) \mid n \in \mathbb{Z}^- \text{ and } m \in \mathbb{N}\} \cup \{(-n, -(m)) \mid n, m \in \mathbb{N}^+ \text{ and } m <_{\mathbb{N}} n\}$$

We define the relation $>_{\mathbb{Z}}$ ([greater than](#)) via $(n > m) :\Leftrightarrow (m < n)$ and the relations $\leq_{\mathbb{Z}}$ and $\leq_{\mathbb{Z}}$ as the [reflexive extensions](#).

less than The $<$ relation is the [transitive](#) closure of the relation $\{(n, s(n)) \mid n \in \mathbb{N}\}$, and \leq its [transitive-reflexive](#) closure. $>$ and \leq are the corresponding [converse relations](#).

For $a < b$ we say that a is [less than](#) b . ¹

EdN:1

library A python [library](#) is a python file with a collection of [functions](#), [classes](#), and [methods](#). It can be loaded via the `import` command.

license A [license](#) is an authorization (by the [licensor](#)) to use the licensed material (by the [licensee](#)).

licensee Defined along with [license](#)

licensor Defined along with [license](#)

line See [text line](#)

line Defined along with [graph](#)

line feed character Defined along with [text line](#)

line number Defined along with [text line](#)

linear ordering We call a [partial ordering](#) R a [linear ordering](#) (or [simple ordering](#) or [total ordering](#)), iff $a \leq b$ or $b \leq a$ for all $a, b \in A$.

linearly ordered set We call a structure $\langle S, \leq \rangle$ of a set S and a [total ordering](#) \leq a [linearly](#) or [totally ordered set](#).

¹EDNOTE: continue for the others

linked Defined along with [path](#)

list A [list](#) or [sequence](#) is an [abstract data type](#) that represents a [finite](#) number of ordered [elements](#), where the same [value](#) may occur more than once.

list See [sequence](#)

list constructor We call `[[seq]]` the [list constructor](#).

literary work Defined along with [copyrightable work](#)

local name Defined along with [XML namespace](#)

logical reasoning **Error: The def does not appear to be inside a definition environment. line 42:54**

logical system **Error: The def does not appear to be inside a definition environment. line 51:5**

loop A [loop](#) is a [control structure](#) that allows to execute certain parts of a [program](#) (the [body](#)) multiple times depending on [conditions](#).

lower bound Defined along with [summation](#)

luminous intensity Defined along with [length](#)

machine instruction Defined along with [computing device](#)

main processor See [central processing unit](#)

mantissa Defined along with [floating point number](#)

mantissa Defined along with [scientific notation](#)

map See [dictionary](#)

markdown cell Defined along with [cell](#)

markup See [document markup](#)

markup code Defined along with [document markup](#)

markup format The [control words](#) and composition rules for a particular kind of [markup](#) system determine a [markup format](#). The [markup format](#) used in an [electronic document](#) is called its [document type](#).

mass Defined along with [length](#)

math idiom Mathematicians use a stylized language that

- uses formulae to represent mathematical objects, e.g. $\int_1^0 x^{3/2} dx$
- uses [math idioms](#) for special situations (e.g. *iff, hence, let...be..., then...*)
- classifies statements by role (e.g. [Definition](#), [Lemma](#), [Theorem](#), [Proof](#), [Example](#))

We call this language [mathematical vernacular](#).

math idiom Defined along with [formulae](#)

mathematical structure A [mathematical structure](#) combines multiple mathematical objects (the [components](#)) into a new object. Structures are usually given as finite enumerations, where the components have names by which they can be referenced.

mathematical vernacular Defined along with [math idiom](#)

mathematical vernacular Defined along with [formulae](#)

maximum Defined along with [minimum](#)

maximum Defined along with [minimum](#)

mebi Defined along with [binary unit prefix](#)

media See [medium](#)

media content Defined along with [medium](#)

medium A [medium](#) (plural [media](#)) is a [communication medium](#) or [storage medium](#). The [information](#) conveyed or [stored](#) is called the [media content](#).

mega Defined along with [prefixes](#)

member Defined along with [set](#)

membership Defined along with [set](#)

memory The [memory](#) (also [primary storage](#)) is a [storage subsystem](#) in a [computer](#) that stores [information](#) for immediate use by its [CPU](#).

message Defined along with [communication medium](#)

metasyntactic variable Defined along with [pseudocode](#)

method Most important [HTTP](#) request [methods](#). (5 more less prominent)

GET	Requests a representation of the specified resource.	safe
PUT	Uploads a representation of the specified resource.	idempotent
DELETE	Deletes the specified resource.	idempotent
POST	Submits data to be processed (e.g., from a web form) to the identified resource.	

method Defined along with [object-oriented programming](#)

micro Defined along with [prefixes](#)

milli Defined along with [prefixes](#)

minimum The [minimum](#) $\min S$ ([maximum](#) $\max S$) of an [ordered set](#) S is that element m (if it exists), such that all other members of S are smaller (larger) than m . We write $\min(a_1, \dots, a_n)$ for $\min\{a_1, \dots, a_n\}$ and $\max(a_1, \dots, a_n)$ for $\max\{a_1, \dots, a_n\}$.

If e is an [expression](#) and φ a condition (in a [variable](#) x), we write $\max_{\varphi}(e)$ for $\max\{e \mid \varphi\}$ and call it the [maximum](#) for e over φ . Analogously, we write $\min_{\varphi}(e)$ for $\min\{e \mid \varphi\}$ and call it the [minimum](#) for e over φ

minimum Defined along with [minimum](#)

minute A [minute](#) is 60 [seconds](#), an [hour](#) is 60 [minutes](#), a [day](#) is 24 [hours](#), and a [week](#) is seven [days](#).

modulo Defined along with [equivalence class](#)

modulus Defined along with [integer division](#)

modulus Defined along with [division](#)

month A [month](#) is between 27 and 31 [day](#), depending on which month of the calendar and year it is.

multi-relation expression A [multi-relation expression](#) is built up from binary relations via conjunction: $aRbSc\dots$ holds, iff $R(a, b)$ holds and also $bSc\dots$

multiplication [Multiplication](#) is extended to \mathbb{Z} by defining the [product](#) by cases:

$$a \cdot b := \begin{cases} |a| \cdot |b| & \text{if } a, b \in \mathbb{N} \text{ or } a, b \in \mathbb{Z}^- \\ -(|a| \cdot |b|) & \text{else} \end{cases}$$

multiplication We define [multiplication](#) on \mathbb{Q} : The [product](#) $\frac{a}{b} \cdot \frac{c}{d}$ is $\frac{a \cdot c}{b \cdot d}$.

multiplication [Multiplication](#) computes the [product](#) $a \cdot b$ (also written as ab or $a \times b$) of [natural numbers](#) a and b . It is defined by the equations $x \cdot 0 = 0$ and $x + s(y) = x + x \cdot y$.

multiplication [Multiplication](#) computes the [product](#) $a \cdot b$ of [numbers](#) a and b .

multiplication Defined along with [absolute value](#)

musical work Defined along with [copyrightable work](#)

mutable The last two items touch a somewhat delicate subject in programming. [Mutable](#) an [immutable data structures](#): the former can be changed in-place – as we have above with the `.set` method, and the latter cannot. Both have their justification and respective advantages. [Immutable data structures](#) are “safe” in the sense that they cannot be changed unexpectedly by another part of the [program](#), they have the disadvantage that every time we want to have a variant, we have to copy the whole object. [Mutable](#) ones do not – we can change in place – but we have to be very careful about who accesses them when.

This is also the reason why we spoke of “dictionary-like interface” to [XML trees](#) in `lxml`: [dictionaries](#) are [immutable](#), while [XML trees](#) are not.

mutually recursive Defined along with [recursive](#)

name Defined along with [subroutine](#)

namespace prefix Defined along with [namespace declaration](#)

namespace declaration Defined along with [XML document tree](#)

namespace declaration [namespace declaration](#) is an attribute `xmlns:⟨prefix⟩=|` whose value is an [XML namespace](#) n on an [XML element](#) e . The first associates the [namespace prefix](#) `⟨prefix⟩` with the [namespace](#) n in e : Then, any [XML element](#) in e with a [prefixed name](#) `⟨prefix⟩:⟨name⟩` has [namespace](#) n and [local name](#) `⟨name⟩`.

A [default namespace declaration](#) `xmlns= d` on an element e gives all elements in e whose name is not [prefixed](#), the [namespace](#) d .

[Namespace declarations](#) on [subtrees](#) shadow the ones on [supertrees](#).

nano Defined along with [prefixes](#)

natural number The set \mathbb{N} of [natural numbers](#) is the set $\{0, 1, 2, \dots\}$. They are constructed by iteration of the [successor function](#) over [zero](#).

navigating Defined along with [hypertext](#)

negative integer	Defined along with non-negative integer
negative rational number	Defined along with positive rational number
negative real number	Defined along with positive real number
node	Defined along with graph
node	Defined along with undirected graph
non-empty	Defined along with empty set
non-negative integer	We use $\mathbb{Z}_0^+ := \mathbb{N}$ and $\mathbb{Z}_0^- := \{x \in \mathbb{Z} \mid x = (-y) \text{ for some } y \in \mathbb{N}\}$ and call them the non-negative integers and non-positive integers respectively. Analogously, we use $\mathbb{Z}^+ := \{x \in \mathbb{Z}_0^+ \mid x \neq 0\}$ and $\mathbb{Z}^- := \{x \in \mathbb{Z}_0^- \mid x \neq 0\}$ for the positive integers and negative integers .
positive rational number	Defined along with positive rational number
positive real number	Defined along with positive real number
non-positive integer	Defined along with non-negative integer
positive rational number	Defined along with positive rational number
non-positive real number	Defined along with positive real number
nondeterministic Turing machine	<p>A nondeterministic Turing machine (NTM) is a septuple $\mathcal{M} := \langle \mathcal{A}, \mathcal{S}, b, \Sigma, s_0, \mathcal{F}, \mathcal{R} \rangle$, where</p> <ul style="list-style-type: none"> – \mathcal{A} is a set called the alphabet, – \mathcal{S} is a set of states, – $b \in \mathcal{A}$ the blank symbol. – $\Sigma \subseteq \mathcal{A}$ the input symbols – $s_0 \in \mathcal{S}$ is the initial state, – $\mathcal{F} \subseteq \mathcal{S}$ is the set of accepting or final states, – and $\mathcal{R} \subseteq ((\mathcal{S} \setminus \mathcal{F}) \times \mathcal{A}) \times (\mathcal{S} \times \mathcal{A} \times \{R, L\})$; it is called the transition relation. <p>\mathcal{M} is called a deterministic Turing machine (DTM) if \mathcal{R} is a function $\mathcal{R}: (\mathcal{S} \setminus \mathcal{F}) \times \mathcal{A} \rightarrow \mathcal{S} \times \mathcal{A} \times \{R, L\}$, then it is called the transition function.</p> <p>If it is irrelevant – or clear from the context – whether \mathcal{M} is deterministic or not, we often just speak of Turing machines without a qualifier.</p>
nonempty string	Let A be an alphabet, then we define the sets $A^+ := \bigcup_{i \in \mathbb{N}^+} A^i$ of nonempty strings and $A^* := A^+ \cup \{\epsilon\}$ of strings .
nonempty word	Let A be an alphabet, then we define the sets $A^+ := \bigcup_{i \in \mathbb{N}^+} A^i$ of nonempty words (nonempty strings) and $A^* := A^+ \cup \{\epsilon\}$ of words (strings).
normalized	Defined along with scientific notation
not fully specified	Defined along with variable
number	A number is a mathematical object used to count or measure other objects.
number system	A number system is simply a set of numbers .
numeral	The representation of a number in a numeral system is called a numeral .

- numeral system** A **numeral system** (or **system of numeration**) is a writing system for expressing **numbers**, that is, a mathematical notation for representing numbers of a given set, using symbols in a consistent manner.
- numerator** Defined along with **rational number**
- object** In python all **values** belong to a **class**, which provide special **functions** we call **methods**. **Values** are also called **objects**, to emphasise **class** aspects. **Method** application is written with **dot notation**: `⟨obj⟩.⟨meth⟩(⟨args⟩)` corresponds to `⟨meth⟩(⟨obj⟩,⟨args⟩)`.
- object** Defined along with **object-oriented programming**
- object-oriented programming** **Object-oriented programming** (OOP) is a **programming** paradigm based on the concept of an **objects**. **Objects** can contain **data**, in the form of **fields** (often known as **attribute** or **properties**) to represent **object** properties. **Object** behavior is specified via **procedures** (called **methods** in OOP).
- octal** Defined along with **unary**
- on** Defined along with **path**
- one** Defined along with **zero**
- one-to-one** See **injective**
- one-to-one correspondence** See **bijjective**
- onto** See **surjective**
- open source** See **Free/Libre/Open-Source Software**
- opened** Once a **file** has been **opened**, the **CPU** can **write** to it and **read** from it. After use a file should be **closed** to protect it from accidental **reads** and **writes**.
- opening tag** For communication this tree is serialized into a balanced bracketing structure, where
- an inner **element node** is represented by the brackets `<el>` (called the **opening tag**) and `</el>` (called the **closing tag**),
 - the **leaves** of the **XML tree** are represented by **empty element tags** (serialized as `<el></el>`, which can be abbreviated as `<el/>`,
 - and text nodes (serialized as a sequence of UniCode characters).
 - An **element node** can be annotated by further information using **attribute nodes** — serialized as an **attribute** in its **opening tag**.
- operating system** An **operating system** (OS) is a system **program** that manages **computer hardware**, **software** resources, and provides common **services** for **computer programs**.
- operator** An **operator** is a **function** that differs syntactically (e.g. by using infix notation) or semantically (in evaluation strategy or argument passing mode) from usual **functions**.
- operator application** Defined along with **expression**
- out-degree** Defined along with **in-degree**
- outdegree** Defined along with **indegree**
- output subsystem** Defined along with **information processing system**

owner There are various legal entities (e.g. persons, states, companies, associations, ...) that can have **ownership** over a **property** p . We call them the **owners** of p .

ownership **Ownership** is the state or fact of exclusive rights and control over **property**, which may be a physical object, land/real estate or intangible object.

pair Defined along with **Cartesian product**

pair Defined along with **set of pairs**

pair set See **unordered pair**

parameter Defined along with **subroutine**

parent Defined along with **tree**

parent Defined along with **tree**

partial function $f \subseteq X \times Y$, is called a **partial function**, iff for all $x \in X$ there is at most one $y \in Y$ with $(x, y) \in f$.

partial function A **relation** $f \subseteq X \times Y$, is called a **partial function** with **domain** X (write **dom**(f)) and **codomain** Y (write **codom**(f)), iff for all $x \in X$ there is at most one $y \in Y$ with $(x, y) \in f$. We write $f: X \rightarrow Y; x \mapsto y$ and $f(x) = y$ instead of $(x, y) \in f$. We say that $f(x)$ is the **application** of f to x and call x the **argument** of f .

partial function space Given sets A and B we will call the set $A \rightarrow B$ ($A \rightarrow B$) of all (partial) functions from A to B the **(partial) function space** from A to B .

partial order See **partial ordering**

partial ordering We call a **preorder** $\leq \subseteq A \times A$ on A a **partial ordering** (or **partial order**), iff it is **antisymmetric**. We associate with \leq a **strict ordering** $< := \{(a, b) \in \leq \mid a \neq b\}$.

We often also use the **converse relations** \geq and $>$.

partially ordered set We call a structure $\langle S, \leq \rangle$ of a set S and a **partial ordering** \leq an **partially ordered set** or **poset**.

patent Defined along with **intellectual property**

path Given a directed graph $G = \langle V, E \rangle$, then we call a vector $p = \langle v_0, \dots, v_n \rangle \in V^{n+1}$ a **path** in G iff $(v_{i-1}, v_i) \in E$ for all $1 \leq i \leq n$, $n > 0$.

– v_0 is called the **start** of p (write **start**(p))

– v_n is called the **end** of p (write **end**(p))

– n is called the **length** of p (write **len**(p))

path Given a **directed graph** $G := \langle V, E \rangle$ we call a $n + 1$ -tuple $p = \langle v_0, \dots, v_n \rangle \in V^{n+1}$ a **path** in G iff $(v_{i-1}, v_i) \in E$ for all $1 \leq i \leq n$ and $n > 0$.

– We say that the v_i are nodes **on** p and that v_0 and v_n are **linked** by p .

– v_0 and v_n are called the **start** and **end** of p (write **start**(p) and **end**(p)), the other v_i are called **inner nodes** of p .

– n is called the **length** of p (write **len**(p)).

– We denote the set of **paths** in G with $\Pi(G)$

path Defined along with [uniform resource identifier](#)

pebi Defined along with [binary unit prefix](#)

personal rights The [copyright](#) is a collection of rights on a [copyrighted](#) work;

- **Personal rights:** the owner of the [copyright](#) may
 - * determine whether and how the work is published (right to publish)
 - * determine whether and how her authorship is acknowledged. (right of attribution)
 - * to object to any distortion, mutilation or other modification of the work, which would be prejudicial to his honor or reputation. (droit de respect)
- **Exploitation rights:** the owner of a [copyright](#) has the exclusive right to do, or authorize to do any of the following:
 - * to reproduce the copyrighted work in copies (or phonorecords);
 - * to prepare derivative works based upon the copyrighted work;
 - * to distribute copies of the work to the public by sale, rental, lease, or lending;
 - * to perform the copyrighted work publicly;
 - * to display the copyrighted work publicly; and
 - * to perform the copyrighted work publicly by means of a digital-audio transmission.

peta Defined along with [prefixes](#)

physical quantity A **physical quantity** is a physical property of a phenomenon, body, or substance, that can be quantified by measurement.

pico Defined along with [prefixes](#)

graphic and sculptural work Defined along with [copyrightable work](#)

pixel Defined along with [raster](#)

place-value notation Defined along with [positional number system](#)

plain text [Digital text](#) is subdivided into **plain text**, where all characters carry the textual information and **formatted text**, which also contains instructions to the [document renderer](#).

point Defined along with [graph](#)

poset See [partially ordered set](#)

positional notation Defined along with [positional number system](#)

positional number system A **positional number system** \mathcal{N} is a pair $\mathcal{N} = \langle D_b, \varphi_b \rangle$ with

- D_b is a finite alphabet of b **digits**. b is called the **base** or **radix** of \mathcal{N}
- assign each digit $d \in D_b$ a number $\varphi_b(d)$ between 0 and $b - 1$.
- Extend φ_b to [sequences](#) of [digits](#) by $\varphi_b(\langle n_k, \dots, n_1 \rangle) := \sum_{i=1}^k \varphi_b(n_i) \cdot b^{i-1}$

positional number system A **positional number system** \mathcal{N} is a triple $\mathcal{N} = \langle D_b, \varphi_b, \psi_b \rangle$ with

- D_b is a finite alphabet of b **digits**. ($b := \#(D_b)$ **base** or **radix** of \mathcal{N})
- $\varphi_b: D_b \rightarrow \{\epsilon, /, \dots, /^{[b-1]}\}$ is bijective (first b unary numbers)
- $\psi_b: D_b^+ \rightarrow \{/\}^*$; $\langle n_k, \dots, n_1 \rangle \mapsto \bigoplus_{i=1}^k \varphi_b(n_i) \odot \exp(/^{[b]}, /^{[i-1]})$ (extends φ_b to string code)

positional number system A **positional number system** is a **numeral system** that uses **positional notation** (also called **place-value notation**) to encode a **number** by a **sequence** of **digits**.

A **positional number system** for a **number system** N is given as a **pair** $P := \langle D, \varphi \rangle$, where D is a **finite** set of **digits** (we call $b := \#(D)$ the **base** or **radix** of P) and a **injective** mapping from D to N .

Positional Notation extends φ to a **bijective** mapping from finite sequences over D to N using the arithmetics of N by interpreting a **finite sequence** a_0, \dots, a_n as a sum of successive powers b^i multiplied by $\varphi(a_i)$. Details vary with N .

positive integer Defined along with **non-negative integer**

positive natural number The set \mathbb{N}^+ of **positive natural numbers** is the set $\{1, 2, 3, \dots\}$.

positive rational number We use $\mathbb{Q}^+ := \{\frac{p}{q} \in \mathbb{Q} \mid p > 0\}$ and $\mathbb{Q}^- := \{\frac{p}{q} \in \mathbb{Q} \mid p < 0\}$ sets of **positive rational numbers** and **negative rational numbers** and $\mathbb{Q}_0^- := \{\frac{p}{q} \in \mathbb{Q} \mid p \leq 0\}$ and $\mathbb{Q}_0^+ := \{\frac{p}{q} \in \mathbb{Q} \mid p \geq 0\}$ sets of **non-positive rational numbers** and **non-negative rational numbers**.

positive real number We use \mathbb{R}^+ and \mathbb{R}^- sets of **positive real numbers** and **negative real numbers** and \mathbb{R}_0^- and \mathbb{R}_0^+ sets of **non-positive real numbers** and **non-negative real numbers**. Here a real number is called positive/negative, iff all rational numbers that approximate it are.

postulate See **axiom**

power Defined along with **exponentiation**

power Defined along with **exponentiation**

power Defined along with **exponentiation**

power Defined along with **exponentiation**

power set the **power set**: $\mathcal{P}(A) := \{S \mid S \subseteq A\}$

power set Let A be a set, then the **power set** $\mathcal{P}(A)$ of A is $\{S \mid S \subseteq A\}$.

pre-image Defined along with **image**

precision Defined along with **floating point number**

predecessor Defined along with **successor**

prefix A string p is called a **prefix** of s (write $p \leq s$), iff there is a string t , such that $s = \text{conc}(p, t)$.
 p is a **proper prefix** of s (write $p \triangleleft s$), iff $t \neq \epsilon$.

prefix Defined along with **subword**

prefix code A (character) code $c: A \rightarrow B^+$ is a **prefix code** iff none of the codewords is a proper prefix to another codeword, i.e.,

$$\forall x, y \in A. x \neq y \Rightarrow (c(x) \not\triangleleft c(y) \wedge c(y) \not\triangleleft c(x))$$

prefixed name Defined along with **namespace declaration**

prefixed unit Defined along with **prefixes**

prefixes The SI system defines a set of 19 **prefixes**, which transform a **unit** into a **prefixed unit** by multiplying it with a power of 10.

Prefix	Symbol	10^n
yotta	Y	10^{24}
zetta	Z	10^{21}
exa	E	10^{18}
peta	P	10^{15}
tera	T	10^{12}
giga	G	10^9
mega	M	10^6
kilo	k	10^3
hecto	h	10^2
deca	da	10^1
deci	d	10^{-1}
centi	c	10^{-2}
milli	m	10^{-3}
micro	μ	10^{-6}
nano	n	10^{-9}
pico	p	10^{-12}
femto	f	10^{-15}
atto	a	10^{-18}
zepto	z	10^{-21}
yocto	y	10^{-24}

preorder We call a binary relation $\leq \subseteq A \times A$ on A a **preorder** (or **quasiorder**), iff it is **reflexive** and **transitive**.

preordered set Defined along with **base set**

presentation MathML Defined along with **Mathematics Markup Language**

primary storage See **memory**

primitive In a **programming language**, a **primitive** is a “basic unit of processing”, i.e. the simplest element that can be given a procedural meaning (its **semantics**) of its own.

procedure Defined along with **subroutine**

process A **process** is an instance of a **program** that is being **executed**.

processor Defined along with **information processing system**

product Defined along with **multiplication**

product Defined along with **multiplication**

product Defined along with **multiplication**

product Defined along with **multiplication**

program Defined along with **programming language**

program Defined along with **Turing machine**

program code See **computer code**

programmer A person involved in **programming** is called a **programmer**.

programming See [computer programming](#)

programming language A [programming language](#) L is a [formal language](#) for specifying sequences [information processing system instructions](#). A [word](#) in L is called a [program](#) of L .

programming paradigm [Programming paradigms](#) are a way to classify [programming languages](#) based on their features. Languages can be classified into multiple [paradigms](#).

projection We call the function $\pi_i: A_1 \times \dots \times A_n \rightarrow A_i; \langle a_1, \dots, a_n \rangle \mapsto a_i$ the (i^{th}) [projection](#).

proof Defined along with [inference](#)

proper prefix Defined along with [prefix](#)

proper prefix Defined along with [subword](#)

proper subgraph Defined along with [subgraph](#)

proper subset A set A is a [proper subset](#) of a set B (written $A \subset B$), iff $A \subseteq B$ but $A \neq B$.

proper substring Defined along with [subword](#)

proper subword Defined along with [subword](#)

proper suffix Defined along with [subword](#)

proper superset A set A is a [proper superset](#) of a set B (written $A \supset B$), iff $B \subset A$.

property Defined along with [object-oriented programming](#)

property Defined along with [ownership](#)

property right [Ownership](#) involves multiple rights (the [property rights](#)), which may be separated and held by different parties.

proset Defined along with [base set](#)

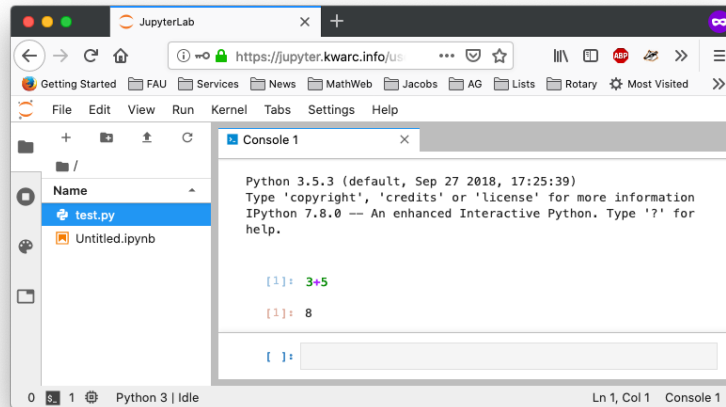
pseudocode [Pseudocode](#) is a plain language description of the steps in an algorithm or another system. [Pseudocode](#) often uses structural conventions of a normal [programming language](#), but is intended for human reading rather than machine reading.

In [pseudocode](#) we often use [metasyntactic variables](#), i.e. is a specific word or set of words identified as a placeholder that is intended to be modified or substituted before real-world usage.

public domain A work is said to be in the [public domain](#), if no [copyright](#) applies, otherwise it is called [copyrighted](#).

punch card A [punch card](#) is a piece of stiff paper that contains digital information represented by the presence or absence of holes in predefined positions.

python console The JupyterLab [python console](#), i.e. a python [interpreter](#) in your [browser](#). (use [this for python interaction and testing](#).)



quantifier Defined along with [formulae](#)

quasiorder See [preorder](#)

query Defined along with [uniform resource identifier](#)

quotient Defined along with [division](#)

quotient Defined along with [division](#)

quotient Defined along with [rational number](#)

quotient set Defined along with [equivalence class](#)

quotient space Defined along with [equivalence class](#)

quotient space Defined along with [equivalence class](#)

radix Defined along with [positional number system](#)

radix Defined along with [positional number system](#)

radix Defined along with [positional number system](#)

range A **range** is a [finite sequence](#) of numbers it can conveniently be constructed by the `range` function: `range(⟨start⟩,⟨stop⟩,⟨step⟩)` constructs a **range** from `⟨start⟩` to `⟨stop⟩` with step size `⟨step⟩`.

range Defined along with [variable](#)

raster We call the grid **raster** and each entry in it **pixel** (from “picture element”).

rational number The set \mathbb{Q} of **rational numbers** is defined as $(\mathbb{Z} \times \mathbb{N}^+)/\sim$, where $((p_1, q_1) \sim (p_2, q_3))$, iff $p_1 q_2 = q_2 q_1$. We call $\frac{n}{m} := (n, m)$ a **fraction** or **quotient**; it denotes the **rational number** $[(n, m)]_{\sim}$. In a **fraction** $\frac{n}{d}$, n is called the **numerator** and d the **denominator**.

raw cell Defined along with [cell](#)

raw string literal Defined along with [string literal](#)

read Defined along with [opened](#)

- read-eval-print loop** A **read-eval-print loop** (REPL), also termed an **interactive toplevel** or **language shell**, is a simple, interactive **user interface** that takes single user inputs (i.e., single **expressions** or **instructions**), **evaluates** or **executes** them, and returns the result to the **user**.
- real number** The set \mathbb{R} of **real numbers** is defined as the completion of \mathbb{Q} .
- real part** Defined along with **complex numbers**
- receiver** Defined along with **communication medium**
- recursion** Defined along with **recursive**
- recursive** We call a problem **recursive**, if its solution depends on solutions to smaller instances of the same problem.
Recursion solves **recursive problems** by (**mutually**) **recursive functions** or **types**.
 We call a set $F = \{f_1, \dots, f_n\}$ of **functions** or **types** **mutually recursive** if they **call** each other in their **bodies**.
 We call a **functions** or **types** f **recursive**, iff $\{f\}$ is **mutually recursive**.
- recursive** Defined along with **recursive**
- refer** Defined along with **reference**
- reference** A **reference** is a **value** that enables a **program** to indirectly access a particular **datum** in the **computer's memory** or in some other **storage** device. The **reference** is said to **refer** to the datum, and accessing the datum is called **dereferencing** the **reference**.
- reflexive** A **relation** $R \subseteq A \times A$ is called
- **reflexive** on A , iff $(a, a) \in R$ for all $a \in A$, and
 - **irreflexive** (or **anti-reflexive**) on A , iff $(a, a) \notin R$ for all $a \in A$.
- reflexive** A relation $R \subseteq A \times A$ is called
- **reflexive** on A , iff $\forall a \in A. (a, a) \in R$
 - **irreflexive** on A , iff $\forall a \in A. (a, a) \notin R$
 - **symmetric** on A , iff $\forall a, b \in A. (a, b) \in R \Rightarrow (b, a) \in R$
 - **asymmetric** on A , iff $\forall a, b \in A. (a, b) \in R \Rightarrow (b, a) \notin R$
 - **antisymmetric** on A , iff $\forall a, b \in A. ((a, b) \in R \wedge (b, a) \in R) \Rightarrow a = b$
 - **transitive** on A , iff $\forall a, b, c \in A. ((a, b) \in R \wedge (b, c) \in R) \Rightarrow (a, c) \in R$
 - **equivalence relation** on A , iff R is **reflexive**, **symmetric**, and **transitive**.
- reflexive extension** A **relation** R on A is **reflexive**, iff $\text{Id}_A \subseteq R$, so we call $R \cup \text{Id}_A$ the **reflexive extension** of R .
- regexp** See **regular expression**
- regular expression** A **regular expression** (also called **regexp**) is a formal expression that specifies a set of **strings**.
- relation** $R \subseteq A \times B$ is a (binary) **relation** between A and B .
 If $A = B$ then R is called a **relation on** A .
- relation** $R \subseteq A \times B$ is a (binary) **relation** between A and B .
- relation on** Defined along with **relation**

relation on If $A = B$ then R is called a **relation on** A .

relative complement See [set difference](#)

remainder Defined along with [integer division](#)

remainder Defined along with [division](#)

renewal provision Defined along with [term provision](#)

representative Defined along with [equivalence class](#)

resource See [system resource](#)

return Defined along with [subroutine](#)

return value Defined along with [subroutine](#)

root Defined along with [tree](#)

root For $b \in \mathbb{Z}$, $r \in \mathbb{Z}$ is a b -th **root** of $a \in \mathbb{Z}$ (we write $\sqrt[b]{a}$), if $r^b = a$. The **square root** \sqrt{a} is written as \sqrt{a} .

root Defined along with [tree](#)

root For $b \in \mathbb{Q}$, $r \in \mathbb{Q}$ is a b -th **root** of $a \in \mathbb{Q}$ (we write $\sqrt[b]{a}$), if $r^b = a$. The **square root** \sqrt{a} is written as \sqrt{a} .

root The n -th **root** $\sqrt[n]{a}$ of a number $a \in \mathbb{N}$ is the $r \in \mathbb{N}$ – if it exists, such that $r^n = a$. The **square root** \sqrt{a} is written as \sqrt{a} .

root The n -th **root** $\sqrt[n]{a}$ of a **number** a is that r – if it exists, such that $r^n = a$. The second **root** is also called the **square root** and is written as \sqrt{a} .

root Defined along with [absolute value](#)

routine See [subroutine](#)

safe We call a [HTTP](#) request **safe**, iff it does not change the state in the web server. (except for [server logs](#), [counters](#),... ; [no side effects](#))

scalar fraction Defined along with [scalar multiple](#)

scalar multiple For a quantity $q \in Q$ and a scalar r we define the **scalar multiple** $r \cdot q$ of r and q to be $r \cdot q := r, su$, if $q = su$ for some $u \in U$.

Similarly, we define the **scalar fraction** q/q' of two quantities q and q' to be that $(r/r') \in \mathbb{R}$, such that if $q = ru$ and $q' = r'u$ for some $u \in U$. Note that both operations are well-defined by the axioms above.

scale Defined along with [floating point number](#)

scheme Defined along with [uniform resource identifier](#)

scientific notation In **scientific notation** all numbers are written in the form of $a \times 10^b$, ($a \in \mathbb{R}$ times 10 raised to the power of $b \in \mathbb{Z}$), a is called the **significand**, **mantissa**, or **coefficient**.

$a \times 10^b$ is called **normalized**, iff $1 \leq |a| < 10$

second Since 1967, the **second** has been defined as exactly the duration of 9,192,631,770 periods of the radiation corresponding to the transition between the two hyperfine levels of the ground state of the caesium-133 atom (at a [temperature](#) of 0°K).

second component Defined along with [first component](#)

secondary storage [Secondary storage](#) (also known as [external memory](#) or [auxiliary storage](#)), refers to any form of [storage device](#) and [media](#) that it is not directly accessible by the [CPU](#).

semantics Defined along with [primitive](#)

sender Defined along with [communication medium](#)

sequence python has more [types](#) that behave just like [lists](#), they are called [sequence types](#).

sequence See [list](#)

sequence A [sequence](#) (also called a [list](#)), $(a_n)_{n \in S}$, is a [function](#) whose domain is a [countable, totally ordered set](#) S (e.g. \mathbb{N} , then we often write (a_i)). If S is [infinite](#), we call any sequence $(a_i)_{i \in S}$ on S an [infinite sequence](#), and [finite sequence](#) otherwise. Then the [cardinality](#) of S is called the [length](#) of $(a_i)_{i \in S}$.

Sequences are written as

- 1, 2, 3, 4 for a concrete finite sequence,
- 1, 2, 3, 4, \dots , 10 for a finite sequence with ellipsis,
- x^1, \dots, x^n and x_1, \dots, x_n for a sequence of upper/lower-indexed variables of length n ,
- 1, 2, 3, 4, \dots for an [infinite sequence](#),
- x^1, x^2, \dots and x_1, x_2, \dots for a sequence of upper/lower-indexed variables of length n ,

Given a sequence $a: S \rightarrow T$, and $a \in S$ we write the i^{th} element of a as S_i .

server A [server](#) is a [program](#) or a [computer](#) that provides functionality – called a [service](#)– for other [programs](#) or [computers](#), called [clients](#).

service Defined along with [server](#)

set A [set](#) is a collection of [elements](#). We can represent [sets](#) by

1. listing the [elements](#) (also called [member](#)) within curly brackets: e.g. $\{a, b, c\}$
2. describing the [elements](#) via selection from another [set](#) S using a property P : $\{x \in S \mid x \text{ has property } P\}$. We use $\{x \mid x \text{ has property } P\}$ as a shorthand, if the set to be selected from is obvious from the context.
3. stating [elementhood](#) (also called [membership](#); written $a \in S$) or not ($b \notin S$) outright.

set comprehension Indeed it is very difficult to define something as foundational as a set. We want sets to be collections of objects, and we want to be as unconstrained as possible as to what their elements can be. But what then to say about them? Cantor's intuition is one attempt to do this, but of course this is not how we want to define concepts in math.

a
AA b
 b

So instead of defining sets, we will directly work with representations of sets. For that we only have to agree on how we can write down sets. Note that with this practice, we introduce a hidden assumption: called [set comprehension](#), i.e. that every set we can write down actually exists. We will see below that we cannot hold this assumption.

set difference [set difference](#): $A \setminus B := \{x \mid x \in A \wedge x \notin B\}$

set difference Let A and B be sets, then the **set difference** $A \setminus B$ of A and B is $\{x \mid x \in A \text{ and } x \notin B\}$.

We also call $A \setminus B$ the **relative complement** or simply **complement** of B in A . If A is clear from the context, we write \bar{B} for the complement of B in A .

set inclusion Defined along with **subset**

set of nonempty strings See **nonempty word**

set of pairs Let A and B be **sets**, then the **set of pairs** $A \times B$ of A and B is defined as $\{(a, b) \mid a \in A, b \in B\}$, we call $(a, b) \in A \times B$ a **pair**. $(a, b) = (c, d)$, iff $a = c$ and $b = d$.

set of strings Defined along with **nonempty word**

set of words Defined along with **nonempty word**

shell A **shell** is a **command-line interface** for accessing the **services** of a **computer's operating system**.

sign Defined along with **floating point number**

significand Defined along with **floating point number**

significand Defined along with **scientific notation**

simple Defined along with **cyclic**

simple Defined along with **cyclic**

simple definition **Error: The def does not appear to be inside a definition environment. line 21:49**

simple ordering See **linear ordering**

sink Defined along with **initial**

sink Defined along with **initial**

size Defined along with **finite**

size The **size** $\#(A)$ of a set A is the number of elements in A .

software Defined along with **computing device**

solar A **solar** or **tropical year** is the time between successive spring or autumn equinoxes, or winter or summer solstices, roughly 365 **days**, 5 **hours**, 48 **minutes**, and 46 **seconds**.

solar calendar A **solar calendar** assigns a **day** to each **solar day** and often aligns **years** with the **solar years**.

solar day A **solar day** is the length of time which elapses between the sun reaching its highest point in the sky two consecutive times. The **solar day** on Earth is roughly 86400s.

sound recording Defined along with **copyrightable work**

source **Compiler**: translates a **program** (the **source**) into another **program** (the **binary**) in a much simpler **programming language** for optimized execution on hardware directly.

source See **initial**

source Defined along with **initial**

source language Defined along with **compiler**

- square root** Defined along with [root](#)
- square root** Defined along with [root](#)
- square root** Defined along with [root](#)
- square root** Defined along with [root](#)
- square root** Defined along with [absolute value](#)
- star operator** The [star operator](#) unpacks a [list](#) into an [argument](#) sequence.
 - start** Defined along with [path](#)
 - start** Defined along with [path](#)
 - state** Defined along with [nondeterministic Turing machine](#)
 - state** Defined along with [stateful](#)
- state register** Defined along with [Turing machine](#)
- stateful** A [system](#) S is described as [stateful](#) if it is designed to remember preceding events. The totality of remembered events is called the [state](#) of S .
- step equation** **Error: The def does not appear to be inside a definition environment. line 75:48**
- storage device** A [storage device](#) is any type of [hardware](#) that [stores](#) (i.e. records with the purpose of later returning) [data](#) in a (fixed or removable) [storage medium](#).
- storage medium** A [storage medium](#) is a physical material that holds ([stores](#)) [information](#).
- storage subsystem** Defined along with [information processing system](#)
 - store** Defined along with [storage device](#)
 - store** Defined along with [storage medium](#)
- stream** Many operating systems use files as a primary computational metaphor, also treating other resources like [files](#). This leads to an abstraction of files called [streams](#), which encompass [files](#) as well as e.g. keyboards, printers, and the screen, which are seen as objects that can be read from (keyboards) and written to (e.g. screens). This practice allows flexible use of [programs](#), e.g. re-directing a the (screen) output of a [program](#) to a [file](#) by simply changing the output [stream](#).
- strict ordering** Defined along with [partial ordering](#)
 - string** Defined along with [integer](#)
 - string** python [strings](#) are sequences of UniCode [characters](#).
 - string** Defined along with [nonempty string](#)
 - string** Defined along with [alphabet](#)
- string literal** python uses [string literals](#), i.e character sequences surrounded by one, two, or three sets of matched single or double quotes for string input. The content can contain [escape sequences](#), i.e. the [escape character](#) backslash followed by a code character for problematic characters:

Seq	Meaning	Seq	Meaning
\\	Backslash (\)	\'	Single quote (')
\"	Double quote (")	\a	Bell (BEL)
\b	Backspace (BS)	\f	Form-feed (FF)
\n	Linefeed (LF)	\r	Carriage Return (CR)
\t	Horizontal Tab (TAB)	\v	Vertical Tab (VT)

In triple-quoted **string literals**, unescaped newlines and quotes are honored, except that three unescaped quotes in a row terminate the literal.

Prefixing a **string literal** with a r or R turns it into a **raw string literal**, in which backslashes have no special meaning.

subgraph Let $G := \langle V, E \rangle$ and $G' := \langle V', E' \rangle$ be two graphs. If $V' \subseteq V$ and $E' \subseteq E$, then G' is a **subgraph** of G , written $G' \subseteq G$. If $G' \subseteq G$ and $G' \neq G$, then G' is a **proper subgraph** of G ; we write $G' \subset G$.

subprogram See **subroutine**

subroutine A **subroutine** (also called **routine** or **subprogram**) is a **program** fragment in a **program** P that performs a specific task, packaged as a unit in so that it can be executed (**called**, or **invoked**) by P .

A **subroutine** p consists of an **identifier** (its **name**), a **sequence** xx_1, \dots, x_n of local **identifiers** called **parameters**, and a **program** fragment (called the **body** of p). The **length** n of x is called the **arity** of p .

When P (the **invoker**) **calls** p , then it supplies a list of **values** (called **arguments**) to p : the **parameters** are replaced by the **arguments**, and the **body** is executed in the context where it is **called**. p may or may not **return values** v to P , the **return values**. If it does, it is called a **function** otherwise a **procedure**.

subset A set A is a **subset** of a set B (written $A \subseteq B$), iff all $x \in A$ are members of B . The relation \subseteq is called **set inclusion**.

substring Let A be an alphabet, then we say that a string $s \in A^*$ is a **substring** of a string $t \in A^*$ (written $s \subseteq t$), iff there are strings $v, w \in A^*$, such that $t = vs w$.

substring See **subword**

subsystem An entity S' in a **system** S that is a **system** itself is called a **subsystem** of S .

subtraction **Subtraction** computes the **difference** $a - b$ of a and b which is defined as $a + (- (b))$.

subtraction The **subtraction** operator computes the **difference** $a - b$ of $a \in \mathbb{Q}$ and $\frac{b}{c} \in \mathbb{Q}$ which is defined as $a + \frac{-(b)}{c}$.

subtraction **Subtraction** – computes the **difference** ab of **natural numbers** a and b . It is defined as is that **natural number** c – if it exists, such that $a + c = b$.

subtraction **Subtraction** – computes the **difference** ab of a and b . It is defined as is that **number** c – if it exists, such that $a + c = b$.

subtraction Defined along with **absolute value**

subtree A **subgraph** of a **tree** that is itself a **tree** is called a **subtree**.

subword Let A be an **alphabet**, then we say that a **word** $s \in A^*$ is a **subword (substring)** of a **word** $t \in A^*$ (written $s \subseteq t$), iff there are **words** $v, w \in A^*$, such that $t = vsw$. If $v \neq \epsilon$ or $w \neq \epsilon$, then we call s a **proper subword (proper substring)** of t and write $s \subset t$.

If $v = \epsilon$, then we call s a **prefix** of t and write $s \leq t$, if additionally $w \neq \epsilon$ we call s a **proper prefix** of t and write $s < t$. Similarly, if $w \neq \epsilon$, then we call s a **suffix** of t , if additionally $v \neq \epsilon$ a **proper suffix** of t .

successor We call a unary natural number the **successor (predecessor)** of another, if it can be constructed by adding (removing) a slash. (successors are created by the **s-rule**)

successor function Defined along with **natural number**

suffix Defined along with **subword**

sum Defined along with **addition**

sum Defined along with **summation**

sum Defined along with **addition**

sum Defined along with **addition**

sum Defined along with **addition**

summation **Summation** is iterated **addition**, we define the **sum** over a sequence a_i by

$$\sum_{i=n}^m a_i := \begin{cases} 0 & \text{if } (n \leq m) \\ a_n + (\sum_{i=(n+1)}^m a_i) & \text{else} \end{cases}$$

The variable i is called the **summation index** and n and m the **lower bound** and **upper bound** of the sum respectively, together they specify the **range of summation**.

There are variant summation operators $\sum_{\varphi} a_i$ and $\sum_{i \in S} a_i$. The first one specifies the range of the summation via a formula φ in i and the second one directly by giving a set S .

summation index Defined along with **summation**

summation range Defined along with **summation**

supergraph If G is a **subgraph** of G' , then we call G' a **supergraph** of G .

superset A set A is a **superset** of a set B (written $A \supseteq B$), iff $B \subseteq A$.

supertree If T is a **subtree** of T' , then we call T' a **supertree** of T .

surjective Defined along with **injective**

surjective A function $f: S \rightarrow T$ is called **surjective** or **onto**, iff for all $y \in T$ there is a $x \in S$ with $f(x) = y$.

symbol A **symbol** is a mark, sign or word that indicates, signifies, or is understood as representing an idea, object, or relationship.

symbol table See **dictionary**

symmetric A relation $R \subseteq A \times A$ is called

– **symmetric** on A , iff $(b, a) \in R$ for all $a, b \in A$ with $(a, b) \in R$.

- **asymmetric** on A , iff $(b, a) \notin R$ for all $a, b \in A$ with $(a, b) \in R$.
- **antisymmetric** on A , iff $(a, b) \in R$ and $(b, a) \in R$ imply $a = b$.

symmetric Defined along with [reflexive](#)

symmetric difference The **symmetric difference** $A\Delta B$ of sets A and B is defined as $(A\setminus B)\cup(B\setminus A)$; it is also written as $A\oplus B$ or $A\ominus B$.

syntax [Programming language syntax](#) describes the surface form of the program: the admissible character sequences. It is also a composition of the [syntax](#) for the [primitives](#).

system A **system** is a group of interacting or interrelated entities that form a unified whole. A [system](#) is delineated by its spatial and temporal boundaries, surrounded and influenced by its environment, described by its structure and purpose and expressed in its functioning.

system of numeration See [numeral system](#)

system resource A **system resource**, or simply **resource**, is any physical or virtual component of limited availability within a [computer system](#) or connected to it.

tag [HTML](#) marks up the structure and appearance of text with **tags** of the form $\langle \text{el} \rangle$ (**begin tag**), $\langle / \text{el} \rangle$ (**end tag**), and $\langle \text{el} / \rangle$ (**empty tag**), where el is one of the following

structure	html, head, body	metadata	title, link, meta
headings	h1, h2, ..., h6	paragraphs	p, br
lists	ul, ol, dl, ..., li	hyperlinks	a
multimedia	img, video, audio	tables	table, th, tr, td, ...
styling	style, div, span	old style	b, u, tt, i, ...
interaction	script	forms	form, input, button
Math	MathML (formulae)	interactive graphics	vector graphics (SVG) and canvas (2D bitmapped)

tape Defined along with [Turing machine](#)

tape See [tape specification](#)

tape specification A **tape specification** (also called **tape**) for a [Turing machine](#) $\mathcal{M} := \langle \mathcal{A}, \mathcal{S}, b, \Sigma, s_0, \mathcal{F}, \mathcal{R} \rangle$ is a sequence $t: \mathbb{N} \rightarrow \mathcal{A}$, such that $t^{-1}(c)$ is **finite** for all $c \in \mathbf{A}$ except b .

target language Defined along with [compiler](#)

tebi Defined along with [binary unit prefix](#)

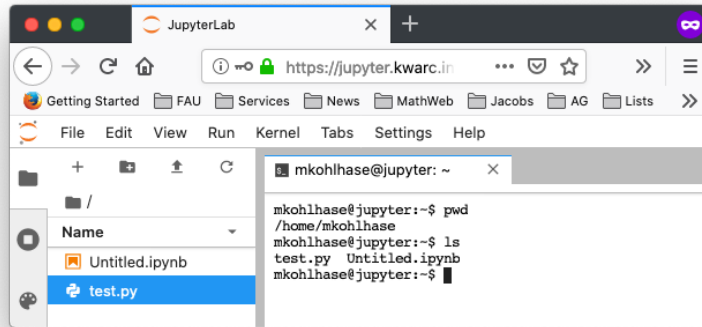
temperature Defined along with [length](#)

tera Defined along with [prefixes](#)

term provision a [license](#) is a regular contract (about [intellectual property](#)) that is handled just like any other contract. (it can stipulate anything the licensor and licensees agree on) in particular a license may

- involve **term**, **territory**, or **renewal** provisions,
- require paying a fee and/or proving a capability, or
- require to keep the licensor informed on a type of activity, and to give them the opportunity to set conditions and limitations.

terminal The JupyterLab **terminal**, i.e. a UNIX **shell** in your browser. (use this for managing files)



terminal Defined along with **initial**

terminal Defined along with **initial**

territory provision Defined along with **term provision**

text editor A **text editor** is a program used for **rendering** and manipulating **text files**.

text file A **text file** is a **file** that is structured as a **sequence** of **encoded characters**. Computer files that are not **text files** are called **binary files**.

text line In practice, **text files** are often processed as a **sequence** of **text lines** (or just **lines**), i.e. substrings separated by the **line feed character** U+000A; LINE FEED (LF). The **line number** is just the position in the sequence.

text node Defined along with **XML document tree**

text node The main remaining functionality in **XML** is the treatment of text. **XML** treats text as special kinds of **node** in the **tree**: **text nodes**. They can be treated just like any other **node** in the **XML tree** in the etree library.

textual content Defined along with **document markup**

theorem A **theorem** is a statement about mathematical objects that we **know to be true**.

time **Time** can be measured by observing a certain number of repetitions of one or another standard cyclical event. Every such event (e.g. the passage of a free-swinging pendulum) constitutes one standard unit.

time Defined along with **length**

total A **relation** $R \subseteq A \times B$ is called **total** iff for all $x \in A$ there is a $y \in B$, such that $(x, y) \in R$.

total $R \subseteq A \times B$ is called **total** iff $\forall x \in A. \exists y \in B. (x, y) \in R$.

total function If $f: X \rightarrow Y$ is a total relation, we call f a **total function** and write $f: X \rightarrow Y$. ($\forall x \in X. \exists^1 y \in Y. (x, y) \in f$)

total function If $f: X \rightarrow Y$ is a **total relation** (i.e. for all $x \in X$ there is a unique $y \in Y$ with $(x, y) \in f$), we call f a **total function** and write $f: X \rightarrow Y$.

- total ordering** See [linear ordering](#)
- totally ordered set** See [linearly ordered set](#)
- trademark** Defined along with [intellectual property](#)
- transition function** Defined along with [nondeterministic Turing machine](#)
- transition relation** Defined along with [nondeterministic Turing machine](#)
- transitive** A relation $R \subseteq A \times A$ is called **transitive** (else **intransitive**) on A , iff $(a, c) \in R$ for all $a, b, c \in A$ with $(a, b) \in R$ and $(b, c) \in R$.
- transitive** Defined along with [reflexive](#)
- transitive closure** Let R be a binary relation, then we call the smallest
- the smallest [transitive relation](#) that contains R the **transitive closure** of R .
 - the smallest [transitive](#) and [reflexive relation](#) that contains R the **transitive-reflexive closure** of R we denote it with R^* .
- transitive-reflexive closure** Defined along with [transitive closure](#)
- transitive-reflexive closure** Let R be a binary relation, then we call the smallest transitive relation that contains R the **transitive-reflexive closure** of R we denote it with R^* .
- tree** A **tree** is a [DAG](#) $G = \langle V, E \rangle$ such that
- There is exactly one initial node $v_r \in V$ (called the **root**)
 - All nodes but the root have in-degree 1.
- We call v the **parent** of w , iff $(v, w) \in E$ (w is a **child** of v). We call a node v a **leaf** of G , iff it is terminal, i.e. if it does not have children.
- tree** A **tree** is a [directed acyclic graph](#) $G := \langle V, E \rangle$ such that
- there is exactly one [initial node](#) $v_r \in V$ (called the **root**), and
 - all [nodes](#) but the [root](#) have [indegree](#) 1.
- We call v the **parent** of w , iff $(v, w) \in E$ (w is a **child** of v). We call a node v a **leaf** of G , iff it is [terminal](#), i.e. if it does not have [children](#). An [ancestor](#) is an iterated [parent](#), and a [descendant](#) an iterated [child](#).
- tropical year** See [solar](#)
- true** Defined along with [truth value](#)
- truth value** There are two **truth values**: **true** (also called **verum**, denoted by T , 1 , or \top) and **false** (or **untrue**, also **falsum**); written as F , 0 , or \perp). The set $\{T, F\}$ of [truth values](#) is denoted with \mathbb{B} .
- twelve** Defined along with [zero](#)
- two** Defined along with [zero](#)
- type** See [data type](#)
- type** Defined along with [variable](#)
- unary** The following positional number systems are in common use.

name	set	base	digits	example
unary	\mathbb{N}_1	1	/	////// ₁
binary	\mathbb{N}_2	2	0,1	0101000111 ₂
octal	\mathbb{N}_8	8	0,1,...,7	63027 ₈
decimal	\mathbb{N}_{10}	10	0,1,...,9	162098 ₁₀ or 162098
hexadecimal	\mathbb{N}_{16}	16	0,1,...,9,A,...,F	FF3A12 ₁₆

unary exponentiation The **unary exponentiation** operation can be defined by the equations $\exp(n, o) = s(o)$ and $\exp(n, s(m)) = n \odot \exp(n, m)$.

unary multiplication The **unary multiplication** operation can be defined by the equations $n \odot o = o$ and $n \odot s(m) = n \oplus n \odot m$.

unary natural numbers We call the representation of natural numbers by slashes on a surface the **unary natural numbers**

unary natural numbers we call these representations **unary natural numbers**.

unary product The **unary product** operation can be defined by the equations $\odot_{i=o}^o n_i = s(o)$ and $\odot_{i=o}^{s(m)} n_i = n_{s(m)} \odot \odot_{i=o}^m n_i$.

unary summation The **unary summation** operation can be defined by the equations $\oplus_{i=o}^o n_i = o$ and $\oplus_{i=o}^{s(m)} n_i = n_{s(m)} \oplus \oplus_{i=o}^m n_i$.

uncountable Defined along with **countable**

undefined at We call a partial function $f: X \rightarrow Y$ **undefined at** $x \in X$, iff $(x, y) \notin f$ for all $y \in Y$. (write $f(x) = \perp$)

undefined at Defined along with **defined at**

undirected edge Defined along with **undirected graph**

undirected graph Defined along with **graph**

undirected graph An **undirected graph** is a pair $\langle V, E \rangle$ such that

- V is a set of **vertices** (or **nodes**) (draw as circles)
- $E \subseteq \{\{v, v'\} \mid v, v' \in V \wedge (v \neq v')\}$ is the set of its **undirected edges** (draw as lines)

unicode Standard Defined along with **universal character set**

unicode standard The **unicode standard** (Unicode) is an industry standard allowing computers to consistently represent and manipulate text expressed in any of the world's writing systems. (currently about 100.000 characters)

uniform resource identifier A **uniform resource identifier (URI)** is a global identifiers of local or network-retrievable documents, or media files (**web resources**). URIs adhere a uniform syntax (grammar) defined in RFC-3986 [BLFM05].

A **URI** is made up of the following **components**:

- a **scheme** that specifies the protocol governing the resource
- an **authority**: the host (authentication there) that provides the resource.
- a **path** in the hierarchically organized resources on the **authority**.
- a **query** in the non-hierarchically organized part of the host data.
- a **fragment identifier** in the resource.

- uniform resource locator** A **uniform resource locator** (**URL**) is a **URI** that that gives access to a **web resource**, by specifying an access method or location. All other **URIs** are called **uniform resource names** (**URN**).
- uniform resource name** Defined along with **uniform resource locator**
- union** Let A and B be sets, then the **union** $A \cup B$ of A and B is defined as $\{x \mid x \in A \text{ or } x \in B\}$.
- union** Let I be a **set** and $\{S_i \mid i \in I\}$ a family of sets, then the **union** $\bigcup_{i \in I} S_i$ over the collection S is $\{x \mid x \in S_i \text{ for some } i \in I\}$.
- union** **union**: $A \cup B := \{x \mid x \in A \vee x \in B\}$
- union over a collection** **union over a collection**: Let I be a **set** and S_i a family of sets indexed by I , then $\bigcup_{i \in I} S_i := \{x \mid \exists i \in I. x \in S_i\}$.
- unit** See **unit of measurement**
- unit of measurement** A **unit of measurement** (or just **unit**) is a definite magnitude of a physical quantity, defined and adopted by convention and/or by law, that is used as a standard for measurement of the same physical quantity. Any other value of the physical quantity can be expressed as a simple multiple of the unit of measurement.
- universal character set** A scalable architecture for representing all the worlds scripts
- The **universal character set** (**UCS**) defined by the ISO/IEC 10646 International Standard, is a standard set of **characters** upon which many character encodings are based.
 - The **unicode Standard** defines a set of standard character encodings, rules for normalization, decomposition, collation, rendering and bidirectional display order
- unknown** Defined along with **variable**
- unordered pair** If A is a set and $x, y \in A$, then the **unordered pair** or **pair set** $\{x, y\}$ is the $p \subseteq A$, such that $z \in p$, iff $z = x$ or $z = y$. The set of all **pair sets** is sometimes denoted with $\binom{A}{2}$.
- untrue** Defined along with **truth value**
- upper bound** Defined along with **summation**
- user** Defined along with **user interface**
- user agent** **HTTP** is used by a client (called **user agent**) to access web resources (addressed by **uniform resource locators** (**URLs**)) via a **HTTP request**. The **web server** answers by supplying the resource (and metadata).
- user interface** A **user interface** (**UI** or simply **interface**), is the means in which a person (the **user**) controls a **software** application or **hardware** device.
- value** Defined along with **dictionary**
- value** Defined along with **variable**
- value** A **value** is the representation of some entity that can be manipulated by a **program**.
- variable** A **variable** is a **memory** location which contains a **value**. It is referenced by an identifier – the **variable name**.

variable A **variable** is a **memory** location which contains a **value**. It is **referenced** by an **identifier** – the **variable name**.

variable A **variable** is an alphabetic character representing a mathematical object, called the **value** of the variable, which is either **arbitrary** (but fixed) or **not fully specified** or unknown – in this case the variable is called an **unknown**. The set of objects a **variable** v can stand for is called the **range** or **type**.

variable assignment A **variable assignment** $\langle\langle\text{var}\rangle\rangle=\langle\langle\text{val}\rangle\rangle$ **assigns** a **value**.

variable binding Defined along with **expression**

variable name Defined along with **variable**

variable name Defined along with **variable**

vector Defined along with **n -dim Cartesian space**

vector Defined along with **n -dimensional Cartesian space**

vector graphics Image representation formats that store shape information instead of individual **pixels**, are referred to as **vector graphics**.

vertex Defined along with **graph**

vertex Defined along with **undirected graph**

verum Defined along with **truth value**

visual markup **Markup** is by no means limited to **visual markup** for documents intended for printing as **?document-markup.ex?** may suggest. There are **aural markup** formats that instruct **document renderers** that transform documents to audio streams of e.g. reading speeds, intonation, and stress.

web IDE A **web IDE** or **cloud IDE**, is a browser-based **integrated development environment**.

web browser A **web browser** is a software application for retrieving (via **HTTP**), presenting, and traversing information resources on the **WWW**, enabling users to view **web pages** and to jump from one page to another.

web page A **web page** is a document (usually marked up in **HTML**) on the **WWW** that can include multimedia data and hyperlinks.

web resource Defined along with **uniform resource identifier**

web server Defined along with **user agent**

web site A **web site** is a collection of related **web pages** usually designed or controlled by the same individual or company.

week Defined along with **minute**

word Defined along with **alphabet**

word processor A **word processor** is a software application, that – apart from being a **document renderer** – also supports the tasks of composition, editing, formatting, printing of **electronic documents**.

work made for hire A **work made for hire** (**WFH**) is a work created by an employee as part of his or her job, or under the explicit guidance or under the terms of a contract.

write Defined along with [opened](#)

year A **year** is either 364 or 365 [day](#) depending whether is is a [leap year](#) or not.

yobi Defined along with [binary unit prefix](#)

yocto Defined along with [prefixes](#)

yotta Defined along with [prefixes](#)

zebi Defined along with [binary unit prefix](#)

zepto Defined along with [prefixes](#)

zero We introduce some abbreviations

- we “abbreviate” o and ‘ ’ by the symbol ‘0’ (called “zero”)
- we abbreviate $s(o)$ and / by the symbol ‘1’ (called “one”)
- we abbreviate $s(s(o))$ and // by the symbol ‘2’ (called “two”)
- ...
- we abbreviate $s(s(s(s(s(s(s(s(s(s(o))))))))))$ and /////////////// by the symbol ‘12’ (called “twelve”)
- ...

zero Defined along with [natural number](#)

zero **Error: The def does not appear to be inside a definition environment. line 40:39**

zetta Defined along with [prefixes](#)

References

[BLFM05] Tim Berners-Lee, Roy T. Fielding, and Larry Masinter. *Uniform Resource Identifier (URI): Generic Syntax*. RFC 3986. Internet Engineering Task Force (IETF), 2005. URL: <http://www.ietf.org/rfc/rfc3986.txt>.

[Hic+14] Ian Hickson et al. *HTML5. A Vocabulary and Associated APIs for HTML and XHTML*. W3C Recommendation. World Wide Web Consortium (W3C), Oct. 28, 2014. URL: <http://www.w3.org/TR/html5/>.