Matriculation Number:

## Name:

# Final Exam General CS (320101)

December 12., 2015

## You have two hours(sharp) for the test;

Write the solutions to the sheet.

The estimated time for solving this exam is 110 minutes, leaving you 10 minutes for revising your exam.

You can reach 110 points if you solve all problems. You will only need 100 points for a perfect score, i.e. 10 points are bonus points.

# Different problems test different skills and knowledge, so do not get stuck on one problem.

	To be used for grading, do not write here													
prob.	1.1	2.1	3.1	3.2	4.1	4.2	5.1	6.1	7.1	7.2	8.1	8.2	Sum	grade
total	5	10	8	5	15	10	15	10	6	8	8	10	110	
reached														

Please consider the following rules; otherwise you may lose points:

- "Prove or refute" means: If you think that the statement is correct, give a formal proof. If not, give a counter-example that makes it fail.
- Always justify your statements. Unless you are explicitly allowed to, do not just answer "yes" or "no", but instead prove your statement or refer to an appropriate definition or theorem from the lecture.
- If you write program code, give comments!

## 1 Easy Points

#### Problem 1.1 (Greek Letters)

Fill in the blanks in the following table of Greek letters. Note that capitalized names 5pt denote capital Greek letters. 3min

Symbol		Ξ	$\eta$			Λ		ω	Ι	
Name	delta			sigma	Psi		Omega	J		chi

#### Solution:

Symbol	δ	[1]	$\eta$	$\sigma$	$\Psi$	Λ	Ω	ω	Ι	$\chi$
Name	delta	Xi	eta	sigma	Psi	Lambda	Omega	omega	Iota	chi

## 2 Induction

#### Problem 2.1 (Proof by Induction)

Consider the following recursively defined function:

- F(0) = 0
- F(1) = 1
- F(n) = F(n-1) + F(n-2) for all  $n \ge 2$ .

Prove by induction or refute that

$$F(n) \le \left(\frac{1+\sqrt{5}}{2}\right)^n$$

Hint:

$$\frac{1+\sqrt{5}}{2} + 1 = \frac{3+\sqrt{5}}{2} = \frac{1+2\sqrt{5}+5}{4} = \left(\frac{1+\sqrt{5}}{2}\right)^2$$

# 3 Relations and Functions

#### Problem 3.1 (Relation Properties)

Let  $A := \{7, 21, 14, 3\}$ , consider the following relations on A:

- $R_1 := \{(14, 14), (21, 3), (3, 14), (21, 14), (14, 3), (3, 21), (14, 21), (7, 7)\}$
- $R_2 := \{(21,3), (7,3), (21,21), (7,7), (7,14), (3,3), (14,14)\}$
- $R_3 := \{(7,7), (21,21), (3,3), (14,14)\}$
- 1. Which of the relations are reflexive, symmetric and transitive?
- 2. Are  $R_1$ ,  $R_2$  and  $R_3$  functions?
- 3. Give example of a linear partial order on A.

10pt 10min

8pt 10min Justify your answers!

## Problem 3.2 (Function Definition)

Let A and B be sets. State the definition of the concept of a partial function with domain A 5pt and codomain B. Also state the definition of a total function with domain A and codomain B. 5min

**Solution**: Let A and B be sets, then a relation  $R \subseteq AB$  is called a **partial/total function**, iff for each  $a \in A$ , there is at most/exactly one  $b \in B$ , such that  $(a, b) \in R$ .

# 4 SML

Problem 4.1 (Substitutions)

Given the following SML datatypes

```
datatype BoolExp = Var of string | True | False | Not of BoolExp |
And of BoolExp*BoolExp | Or of BoolExp*BoolExp;
```

#### **datatype** Substitution = Subst **of** string \* BoolExp;

Write an SML function substitute that takes a list of substitutions and a boolean expression and applies the substitutions to the boolean expression (in the same order as they are listed).

Example:

```
substitute(Subst("x1", True)::Subst("x2", Not(Var("x3")))::nil,
And(Or(Not(Var("x1")), Var("x2")), Not(And(Var("x3"), Var("x1")))))
= And(Or(Not(True), Not(Var("x3"))), Not(And(Var("x3"), True)));
```

#### Solution:

fun subst(Subst(str, expr), Var(v)) = if str=v then expr else Var(v)
 | subst(\_, True) = True
 | subst(\_, False) = False
 | subst(s, Not(a)) = Not(subst(s, a))
 | subst(s, And(a, b)) = And(subst(s, a), subst(s, b))
 | subst(s, Or(a, b)) = Or(subst(s, a), subst(s, b));

fun substitute(nil, expr) = expr
| substitute(h::t, expr) = substitute(t, subst(h, expr));

## Problem 4.2 (Duplicates)

Write an SML function that removes all duplicate elements from a list. For instance10ptremove duplicates([true, true, false]) = [true, false];10min

remove\_duplicates([5,3,12,3,3,2]) = [5,3,12,2];

**Hint:** Write a helper function that removes duplicates, but remembers what it has already found in an argument.

#### Solution:

15pt 15min fun member(a, h::t) = a=h orelse member(a, t)
 | member(\_, nil) = false;
fun helper(found\_already, h::t) =
 if member(h, found\_already) then
 helper(found\_already, t)
 else
h::helper(h::found\_already, t)
 | helper(\_, nil) = nil;
fun remove\_duplicates(l) = helper(nil, l);

# 5 Abstract Datatypes

## Problem 5.1 (Proper Binary Trees)

- 1. Design an abstract data type for proper binary trees storing unary natural numbers. In a proper binary tree, every node other than the leaves has exactly two children.
- 2. Give the representation of the binary tree on the right in your ADT.
- 3. A binary tree has the min-heap property if the value of each node is greater than or equal to the value of its parent. Create an abstract procedure that checks if a proper binary tree has the min-heap property.

**Note:** You may define as many helper procedures as you need. You cannot assume that binary numbers are defined already.

# 6 Formal Languages

#### Problem 6.1 (Codes)

Let  $A = \{0, 1, 2\}$  and  $B = \{0, 1\}$  be alphabets. 1. Specify a prefix code  $p : A \to B^+$ 2. Apply the extension of p to the string 02201. 3. Specify a character code  $c : A \to B^+$  such that (a) c is not a prefix code (b) The extension of c is a string code, i.e. it is injective. 4. Apply the extension of c to the string 02201. 10pt 10

 $\mathbf{2}$ 

3

3

5

# 7 Boolean Algebra

#### Problem 7.1 (Boolean Expression)

Given the Boolean expression

6pt 6min

15pt

15min

 $x_1 * (x_2 + x_3 * x_0) + x_1 * (x_0 * x_3)$ 

1. What are the cost and depth of the expression?

- 2. Find an equivalent expression with smaller cost. What is the new cost?
- 3. Find an equivalent expression with smaller depth. What is the new depth?

#### Solution:

## Problem 7.2 (QMC Algorithm)

Execute the Quine-McCluskey algorithm to get the minimum polynomial for the Boolean 8pt function given by 8min

8min

$x_1$	$x_2$	$x_3$	$x_4$	f
Т	Т	Т	Т	Т
Т	Т	Т	F	F
Т	Т	F	Т	Т
Т	Т	F	F	F
Т	F	Т	Т	Т
Т	F	Т	F	F
Т	F	F	Т	Т
Т	F	F	F	F
F	Т	Т	Т	Т
F	Т	Т	F	F
F	Т	F	Т	F
F	Т	F	F	F
F	F	Т	Т	Т
F	F	Т	F	F
F	F	F	Т	F
F	F	F	F	F

## Solution:

 $QMC_1$ :

$x_1$	$x_2$	$x_3$	$x_4$
Т	Т	Т	Т
T	Т	F	Т
T	F	Т	Т
T	F	F	Т
F	Т	Т	Т
F	F	Т	Т

$x_1$	$x_2$	$x_3$	$x_4$
Т	Т	X	Т
Т	X	Т	Т
X	Т	Т	Т
Т	X	F	Т
Т	F	X	Т
X	F	Т	Т
F	X	Т	Т
$x_1$	$x_2$	$x_3$	$x_4$
Т	X	X	Т
X	X	Т	Т

Therefore the prime implicants are  $x_1 x_4$  and  $x_3 x_4$ 

 $QMC_2$ :

	TTTT	TTFT	TFTT	TFFT	FTTT	FFTT
$x_1 x_4$	Т	Т	Т	Т	F	F
$x_3 x_4$	Т	F	Т	F	Т	Т

Therefore both prime implicants are essential.

Final result:  $f = x_1 x_4 + x_3 x_4$ 

# 8 Propositional Logic

#### Problem 8.1 (Hilbert Axioms)

Proove the $K$ and $S$ axioms of the Hilbert calculus using the tableau method.	$8 \mathrm{pt}$
Recall the definitions of $K$ and $S$ :	8min
1. $K := P \Rightarrow Q \Rightarrow P$	011111

2. 
$$S := (P \Rightarrow Q \Rightarrow R) \Rightarrow (P \Rightarrow Q) \Rightarrow P \Rightarrow R$$

Hint: You can use the derived rules

#### Problem 8.2 (CopyLeft)

Briefly state the copyleft clause in the GNU Public License or in the Creative Commons 10pt licenses, and explain how it works. 10min

**Solution**: The copyleft clause states that if a derived work of a licensed work is distributed, then it has to be licensed in exactly the same license as the licensed work.

This makes sure that anybody who wants to make a derived work of the licensed work, they have to decide whether they

- want to distribute it then they have to license it under the copyleft, and contribute to the Open Source Domain, or
- don't, then they do not have to license it at all (but do not get the benefits of distribution/sale).