Assignment12 – ALC

Problem 12.1 (ALC)

- Consider the following description logic signature
- concept symbols: i (for instructor), s (for student), c (for course), p (for program)
- role symbol m (for is-member-of) used for
 - *instructors* giving a *course*
 - *students* taking a *course*
 - *students* being enrolled in a *degree program*
 - courses being part of a degree program

We use an extension of ALC, in which there are dual roles: there is a role m^{-1} that captures the relation has-as-member, e.g., MK m AI iff $AI m^{-1} MK$.

1. For the *signature* above, give a *concept axiom* that captures that instructors can only be members of *courses*.

Solution: $i \sqsubseteq \forall m.c$

2. Give a *concept axiom* for the above *signature* that captures: *courses* that are taken by a *student*, must be given by an *instructor*.

Solution: $c \sqcap \exists m^{-1}.s \sqsubseteq \exists m^{-1}.i$

3. Calculate the translation to *first-order logic* of $s \sqsubseteq \forall m. \exists m. p$.

Solution: $\forall x.s(x) \Rightarrow (\forall y.m(x,y) \Rightarrow (\exists z.m(y,z) \land p(z)))$

 Given a *first-order model* ⟨D, J⟩, define an appropriate case of the *interpreta*tion mapping for the formula ∀r⁻¹.C.

Solution: $\llbracket \forall r^{-1}.C \rrbracket = \{ u \in D | \text{ for } v \in \mathcal{D}, \text{ if } (v, u) \in \llbracket r \rrbracket \}$, then $v \in \llbracket C \rrbracket \}$

Problem 12.2 (ALC Semantics)

Consider the ALC concepts $\forall R.(C \sqcap D)$ and $\forall R.C \sqcap \forall R.D$.

1. By applying the semantics of \mathcal{ALC} , show that the two are equivalent.

Solution: We have:

$$\llbracket \forall R.(C \sqcap D) \rrbracket$$

= {x \in \mathcal{D}| for all y \in \mathcal{D}, if (x, y) \in [[R]], then y \in [[C \ldot D]]}
= {x \in \mathcal{D}| for all y \in \mathcal{D}, if (x, y) \in [[R]], then y \in ([[C]] \circ [[D]])}

$$\begin{split} \llbracket \forall R.C \sqcap \forall R.D \rrbracket \\ &= \llbracket \forall R.C \rrbracket \cap \llbracket \forall R.D \rrbracket \\ &= \{x \in \mathcal{D} | \text{ for all } y \in \mathcal{D}, \text{ if } (x, y) \in \llbracket R \rrbracket, \text{ then } y \in (\llbracket C \rrbracket \cap Q) \} \end{split}$$

where

$$Q = \{x \in \mathcal{D} | \text{ for all } y \in \mathcal{D}, \text{ if } (x, y) \in [[R]], \text{ then } y \in [[D]] \}$$

Now to prove that sets are equal, consider an $x \in \mathcal{D}$ and see that both conditions are equivalent to

for all $y \in \mathcal{D}$, if $(x, y) \in [[R]]$, then $y \in [[C]]$ and $y \in [[D]]$

2. Translate both formulas to first-order logic and state which FOL formula we would need to prove (e.g., with the ND calculus) to show that the two are equivalent.

Solution: The translation yields

$$C_1(x) = \forall y. R(x, y) \Rightarrow C(y) \land D(y)$$

 $C_2(x) = (\forall y.R(x, y) \Rightarrow C(y)) \land (\forall y.R(x, y) \Rightarrow D(y))$

We need to show

$$\forall x.C_1(x) \Leftrightarrow C_2(x)$$

Problem 12.3 (ALC TBox)

Consider ALC with the following

- primitive concepts: woman, man
- roles: has_child, has_parent, has_sibling, has_spouse

Give an *ALC TBox* that defines the *concepts* person, parent, mother, father, grandmother, aunt, uncle, sister, brother, onlychild, cousin, nephew, niece, fatherinlaw, motherinlaw.

Solution: person = man ⊔ woman parent = person∃has_child.person mother = womanparent father = manparent grandmother = woman∃has_child.parent aunt = woman∃has_sibling.parent uncle = man∃has_sibling.parent sister = woman∃has_sibling.person brother = man∃has_sibling.person onlychild = person∃has_parent.∃has_sibling.parent nephew = man∃has_parent.∃has_sibling.person niece = woman∃has_parent.∃has_sibling.person fatherinlaw = man∃has_child.∃has_spouse.person motherinlaw = woman∃has_child.∃has_spouse.person